

Algorithms and Data Structures

CS-CO-412

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Simple Searching & Sorting Algorithms

Lecture 4

Topic Overview

- Linear Search
- Binary Search
- Bubble Sort
- Selection Sort
- Quick Sort

Linear (Sequential) Search

- Linear (Sequential) Search
- Begin at the beginning of the list
- Proceed through the list, sequentially and element by element,
- Until the key is encountered
or
Until the end of the list is reached

Linear (Sequential) Search

- Note: we treat a list as a general concept, decoupled from its implementation
- The order of complexity is $O(n)$
- The list does not have to be in sorted order

Implementation of linear search in C

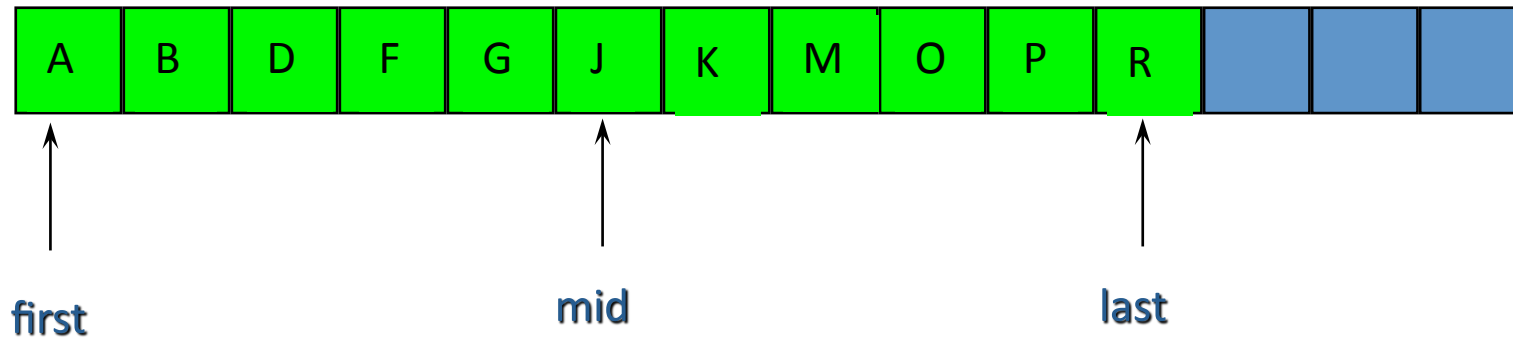
```
int linear_search(item_type s[], item_type key, int low, int high) {  
  
    int i;  
  
    i = low;  
  
    while ((s[i] != key) && (i < high)) {  
        i = i+1;  
    }  
  
    if (s[i] == key) {  
        return (i);  
    }  
    else {  
        return(-1);  
    }  
}
```

Binary Search

- If the list is sorted, we can use a more efficient $O(\log_2(n))$ search strategy
- Check to see whether the key is
 - equal to
 - less than
 - greater than

the middle element

Binary Search



Binary Search

- If key is equal to the middle element, then terminate (found)
- If key is less than the middle element, then search the left half
- If key is greater than the middle element, then search the right half
- Continue until either
 - the key is found or
 - there are no more elements to search

Implementation of Binary_Search

Pseudo-code

```
binary_search(list, key, lower_bound, upper_bound)
```

identify sublist to be searched by setting bounds on search

REPEAT

 get middle element of list

 if middle element < key

 then reset bounds to make the right sublist
 the list to be searched

 else reset bounds to make the left sublist
 the list to be searched

UNTIL list is empty or key is found

Implementation of binary search in C (iterative approach)

```
typedef char item_type;

int binary_search(item_type s[], item_type key, int low, int high) {

    int first, last, mid;

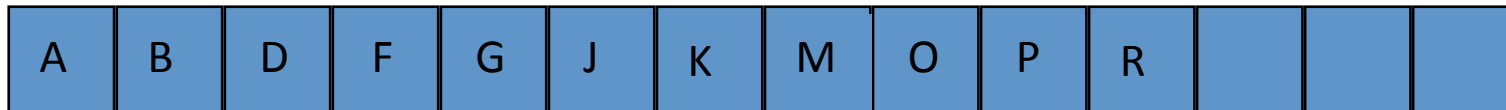
    first = low;
    last  = high;

    do {
        mid = (first + last) / 2;
        if (s[mid] < key) {
            first = mid + 1;
        }
        else {
            last = mid - 1;
        }
    } while ( (first <= last) && (s[mid] != key) );

    if (s[mid] == key)
        return (mid);
    else
        return (-1);

}
```

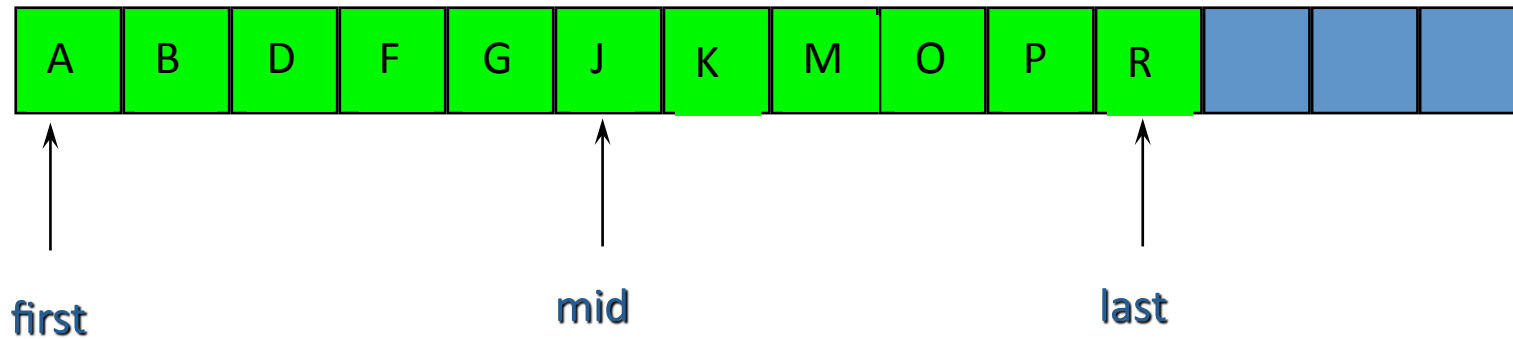
Binary Search



```
first:
last:
mid:
list[mid]:
key:      P
```

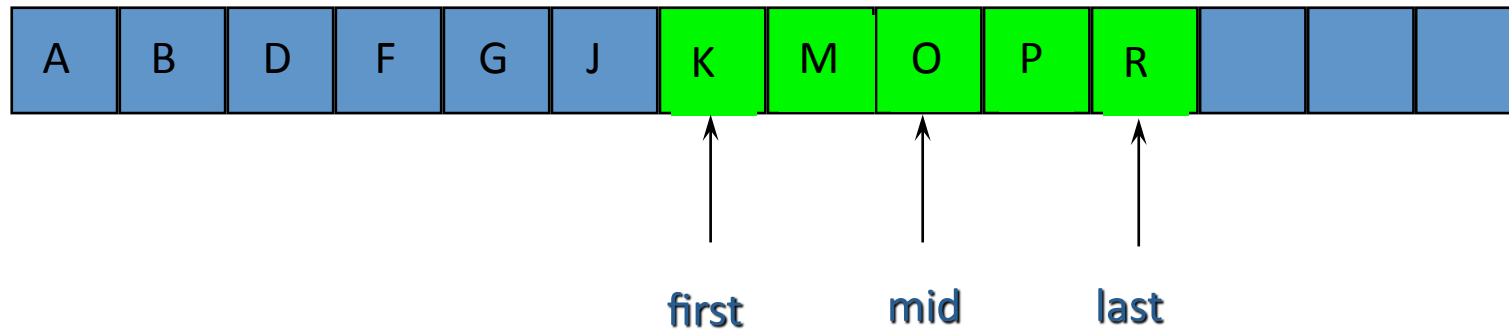
↑ ↑ ↑
first mid last

Binary Search



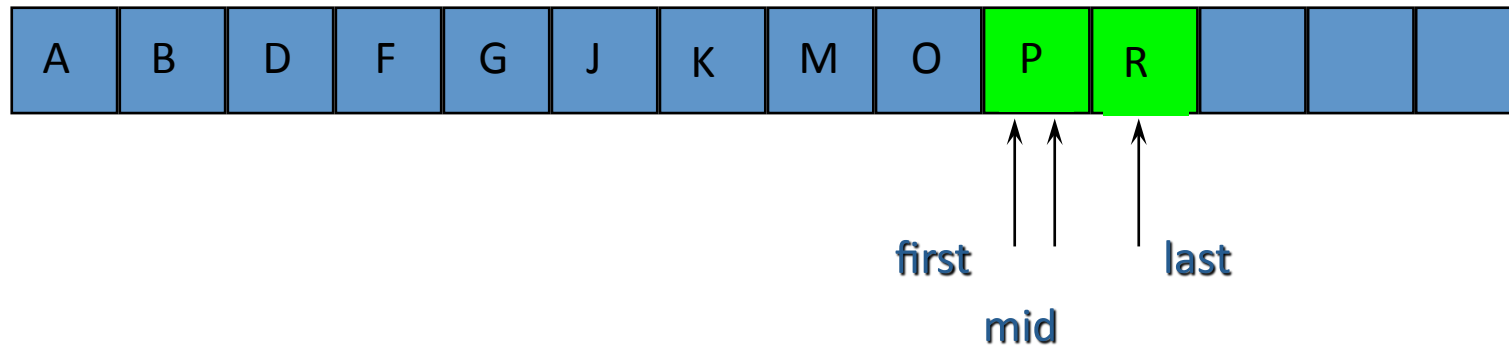
```
first:    1
last:     11
mid:      6
list[mid]: J
key:      P
```

Binary Search



```
first:      1      7
last:      11     11
mid:        6      9
list[mid]:  J      O
key:       P      P
```

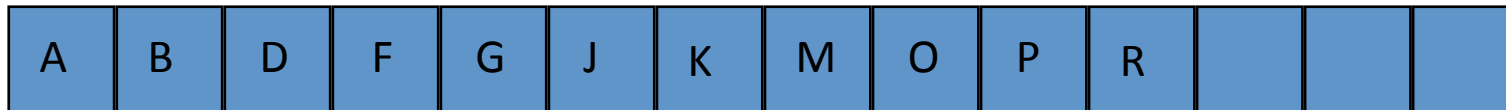
Binary Search



first:	1	7	10
last:	11	11	11
mid:	6	9	10
list[mid]:	J	O	P
key:	P	P	P

← **FOUND!**

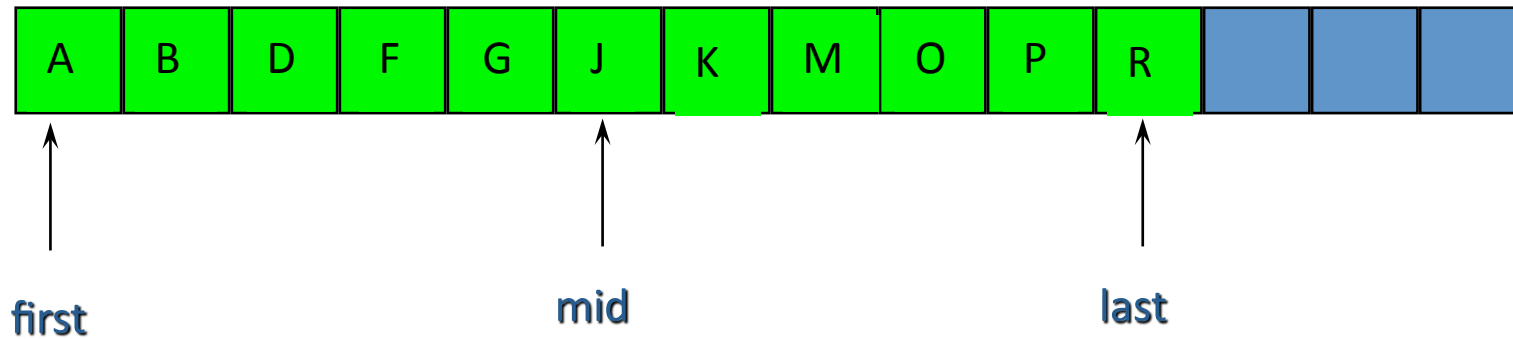
Binary Search



```
first:
last:
mid:
list[mid]:
key:      E
```

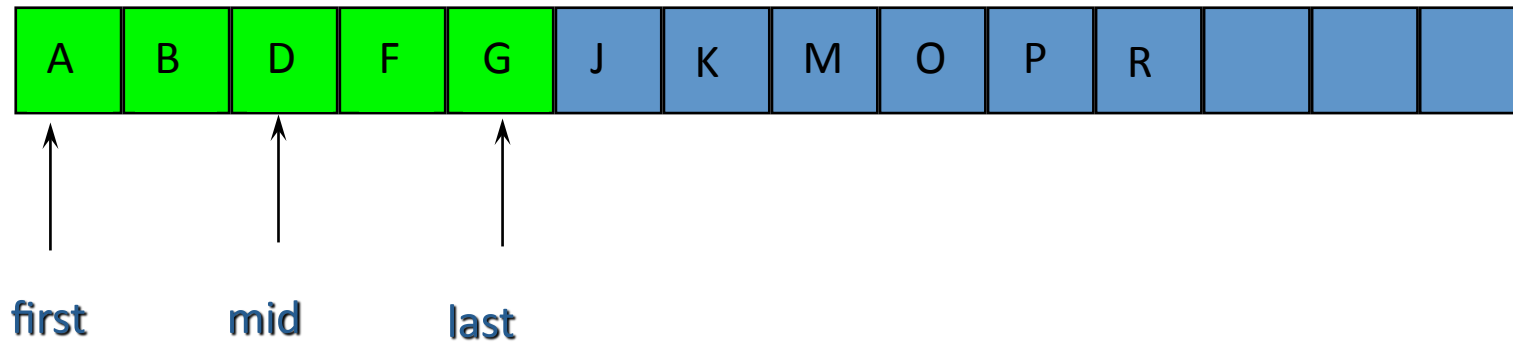
↑ ↑ ↑
first mid last

Binary Search



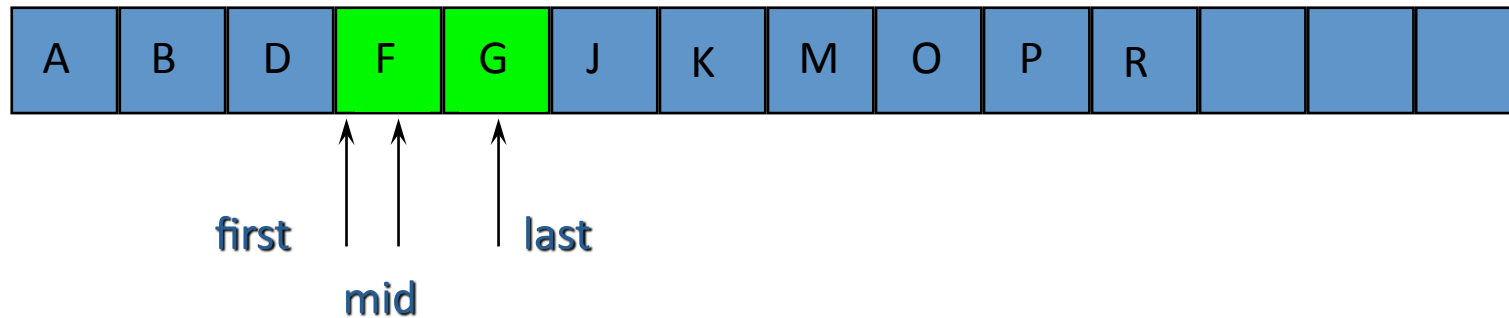
```
first:    1
last:     11
mid:      6
list[mid]: J
key:      E
```

Binary Search



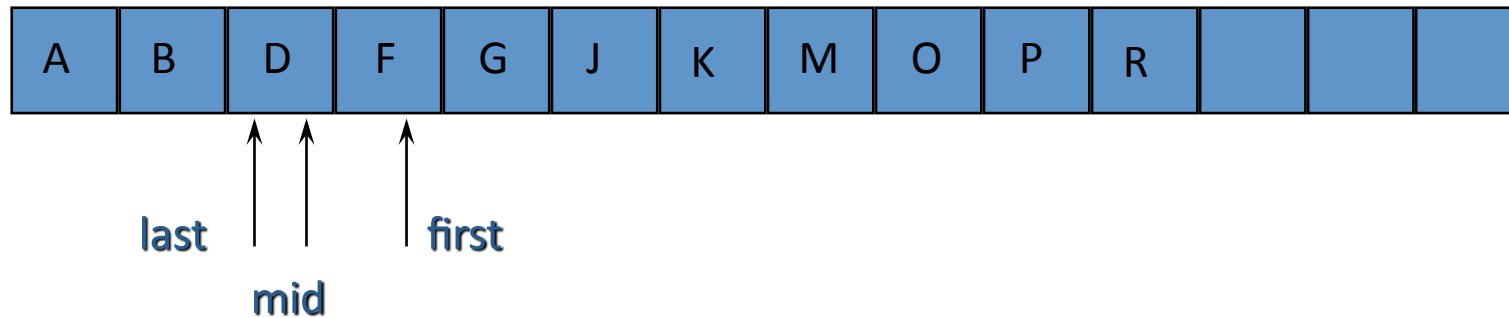
```
first:      1      1
last:      11     5
mid:        6      3
list[mid]: J      D
key:       E      E
```

Binary Search



first:	1	1	4
last:	11	5	5
mid:	6	3	4
list[mid]:	J	D	F
key:	E	E	E

Binary Search



first:	1	1	4	4
last:	11	5	5	3
mid:	6	3	4	3
list[mid]:	J	D	F	D
key:	E	E	E	E

← first > last: NOT FOUND!

Implementation of binary search in C (recursive approach)

```
typedef char item_type;

int binary_search(item_type s[], item_type key, int low, int high) {

    int mid;

    if (low > high) return (-1); /* key not found */

    mid = (low + high) / 2;

    if (s[mid] == key) return(mid);

    if (s[mid] > key) {
        return(binary_search(s, key, low, mid-1));
    }
    else {
        return(binary_search(s, key, mid+1, high));
    }
}
```

Sorting Algorithms

Sorting Algorithms

- Bubble Sort
- Selection Sort
- Quick Sort

Bubble Sort

- Assume we are sorting a list represented by an array A of n integer elements
- Bubble sort algorithm in pseudo-code

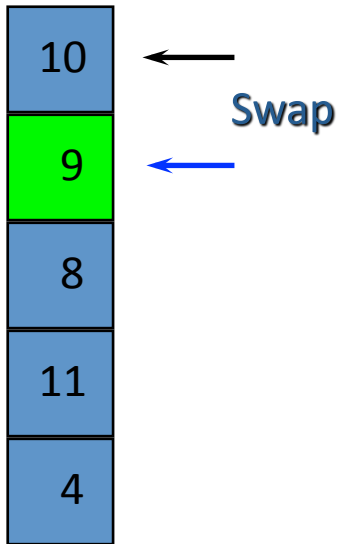
```
FOR every element in the list,  
    proceeding for the first to the last
```

```
    WHILE list element > previous list element  
        bubble element back (up) the list  
        by successive swapping with  
        the element just above/prior it
```

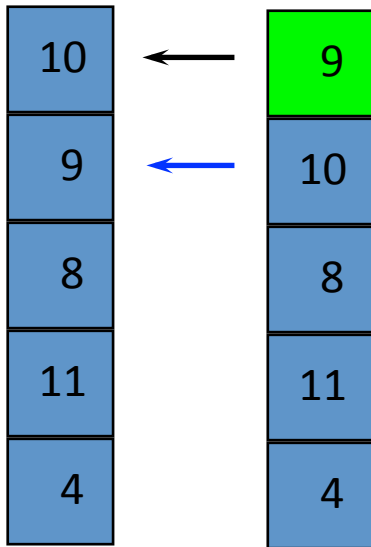
Bubble Sort

10	9	8	11	4
----	---	---	----	---

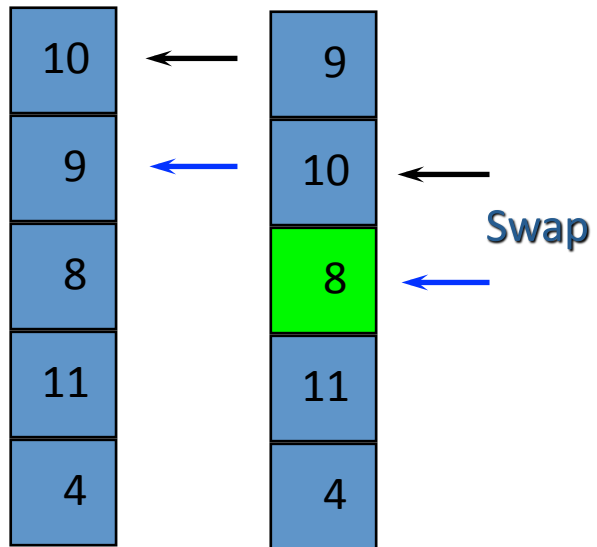
Bubble Sort



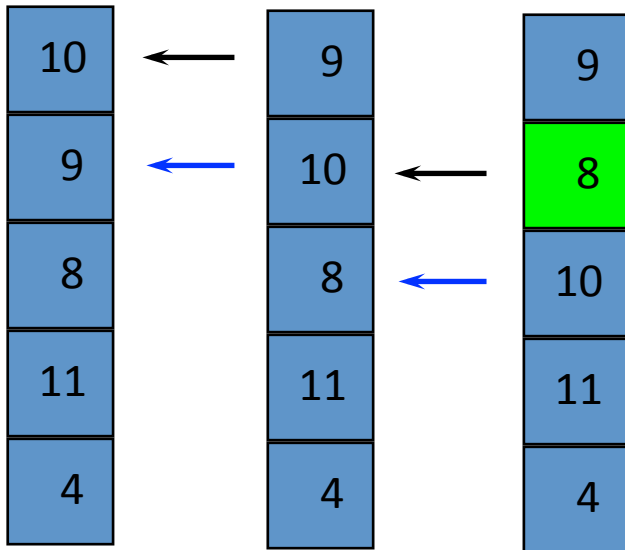
Bubble Sort



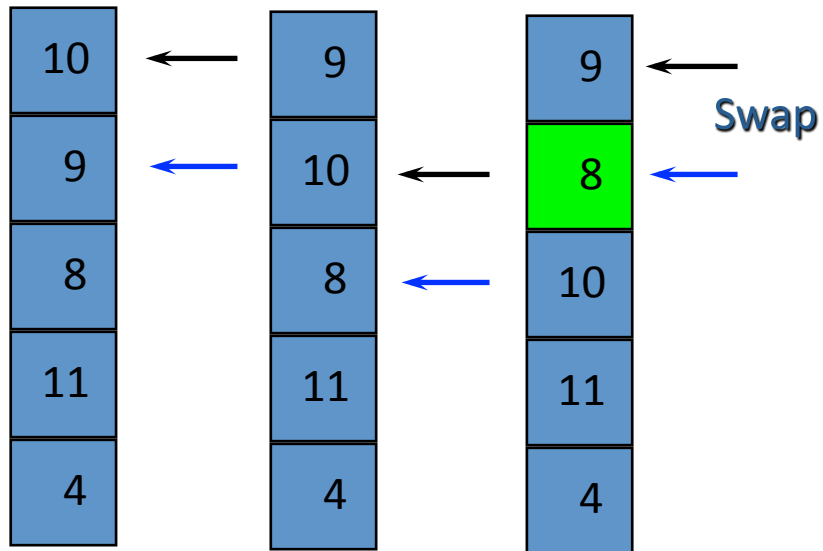
Bubble Sort



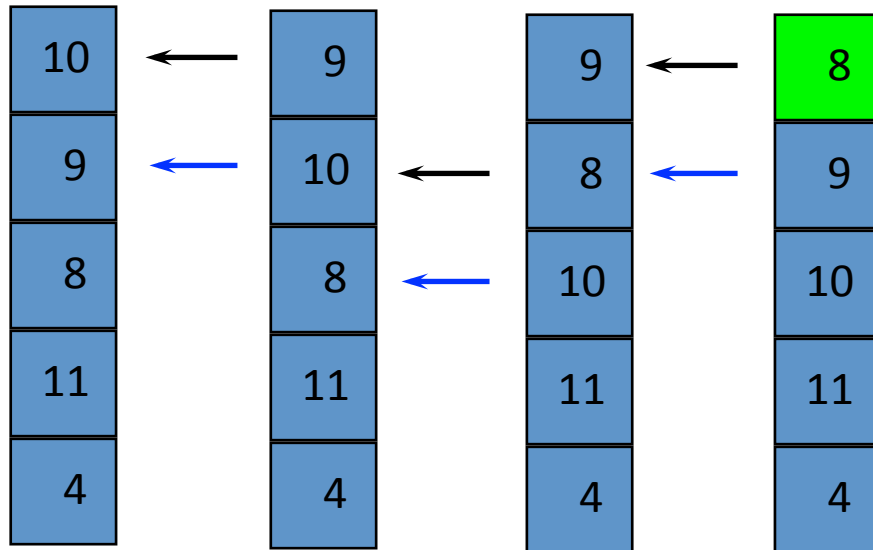
Bubble Sort



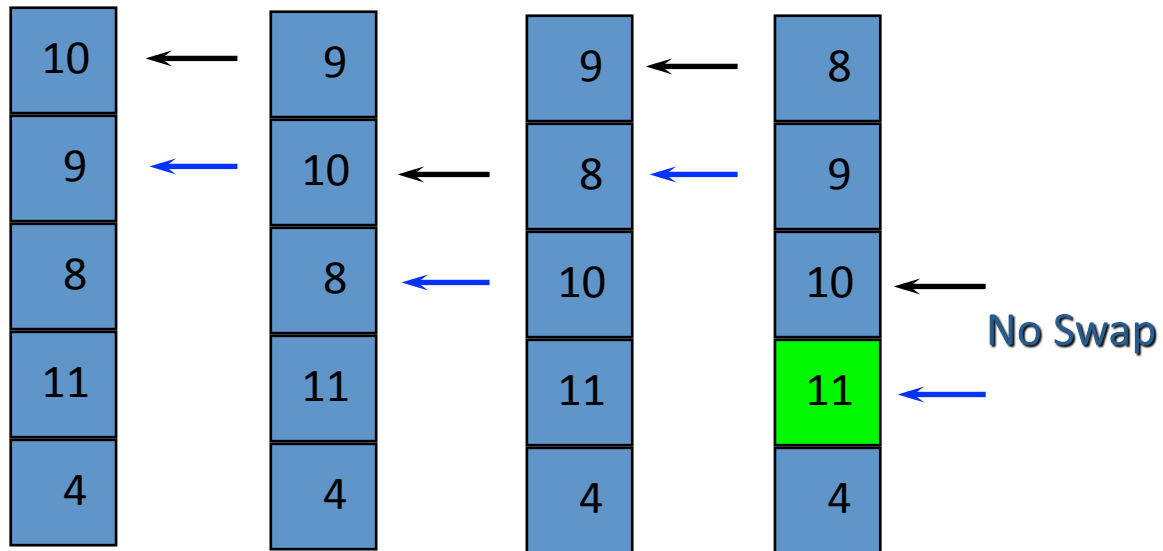
Bubble Sort



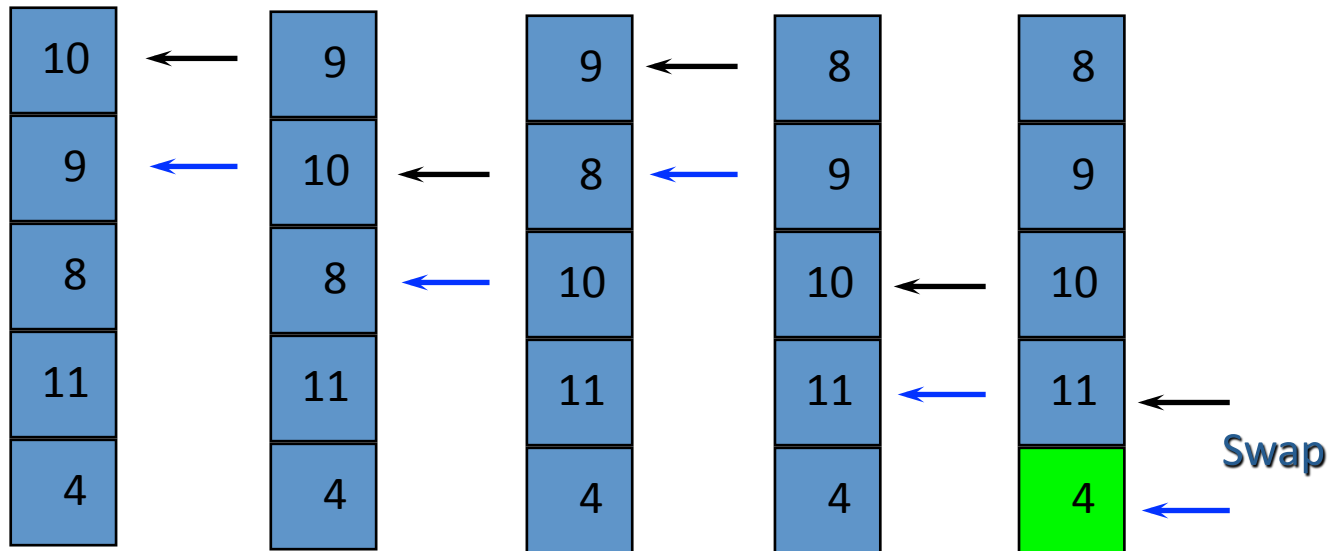
Bubble Sort



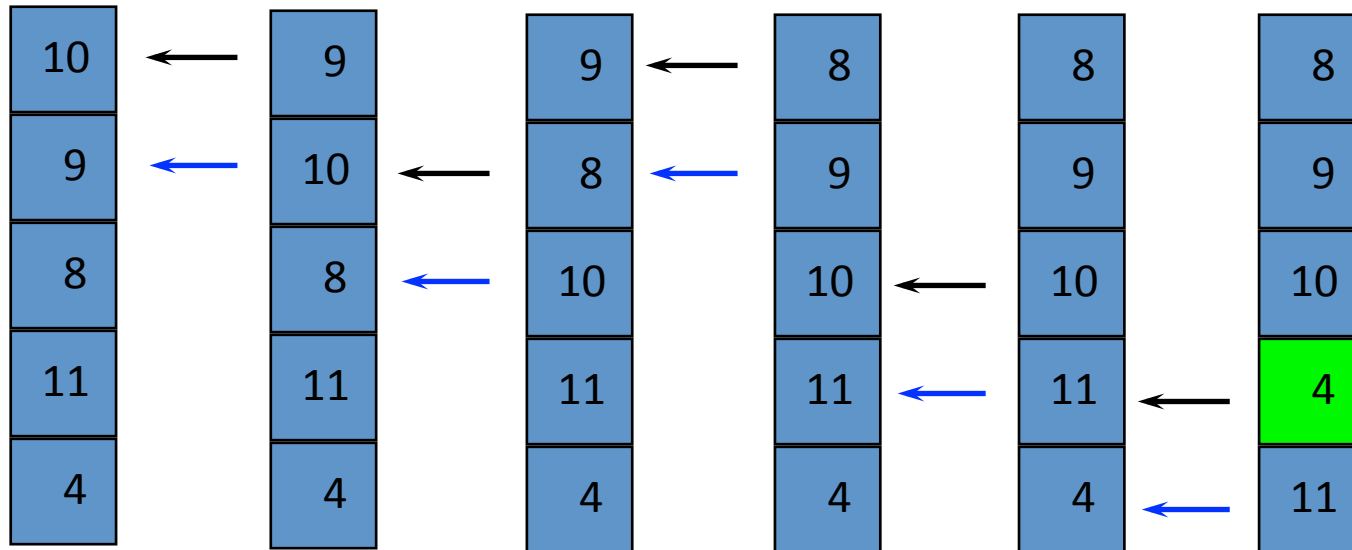
Bubble Sort



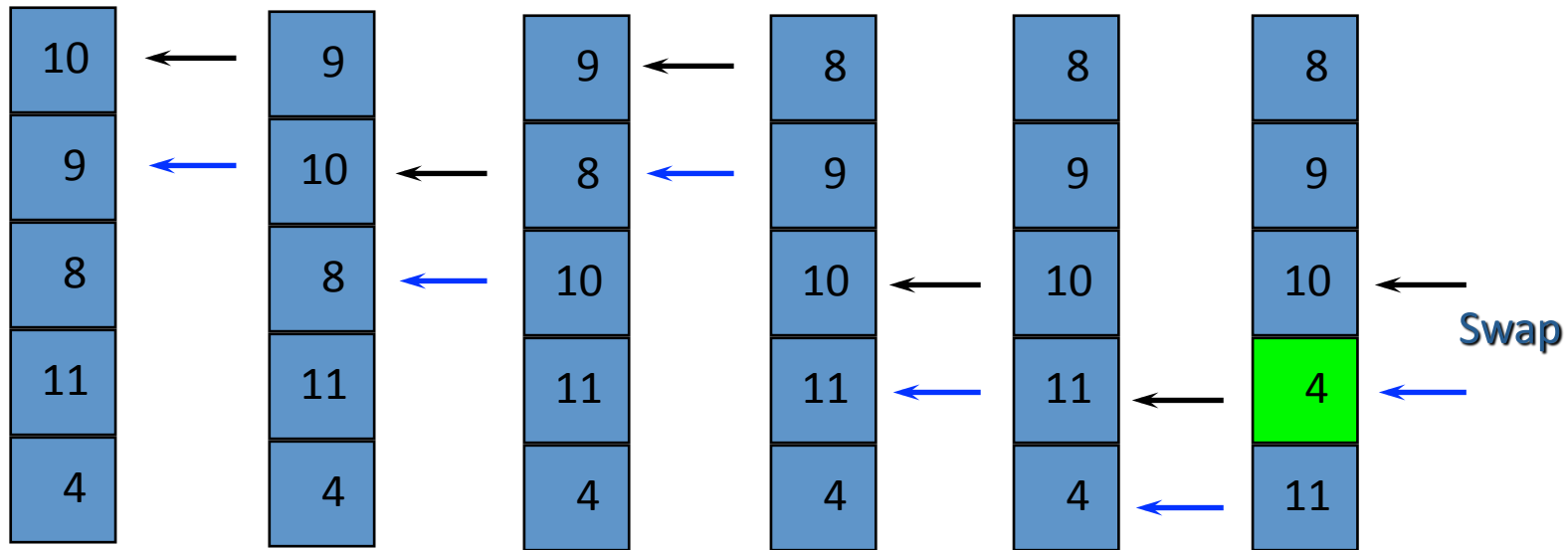
Bubble Sort



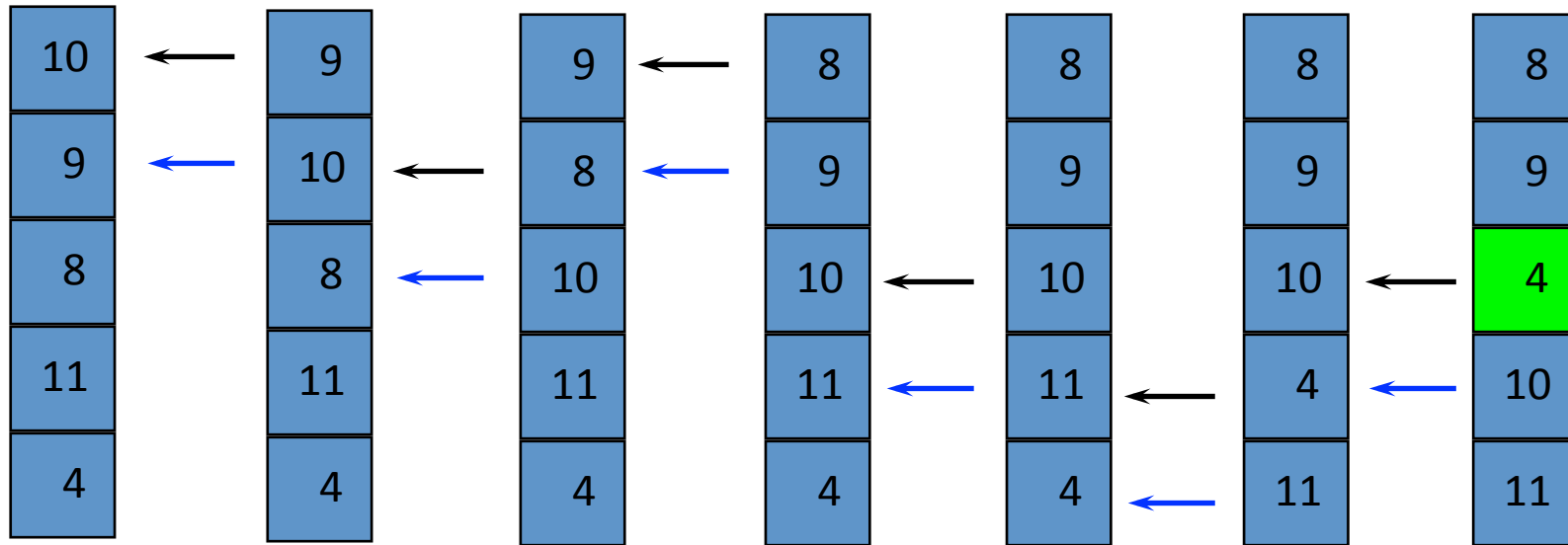
Bubble Sort



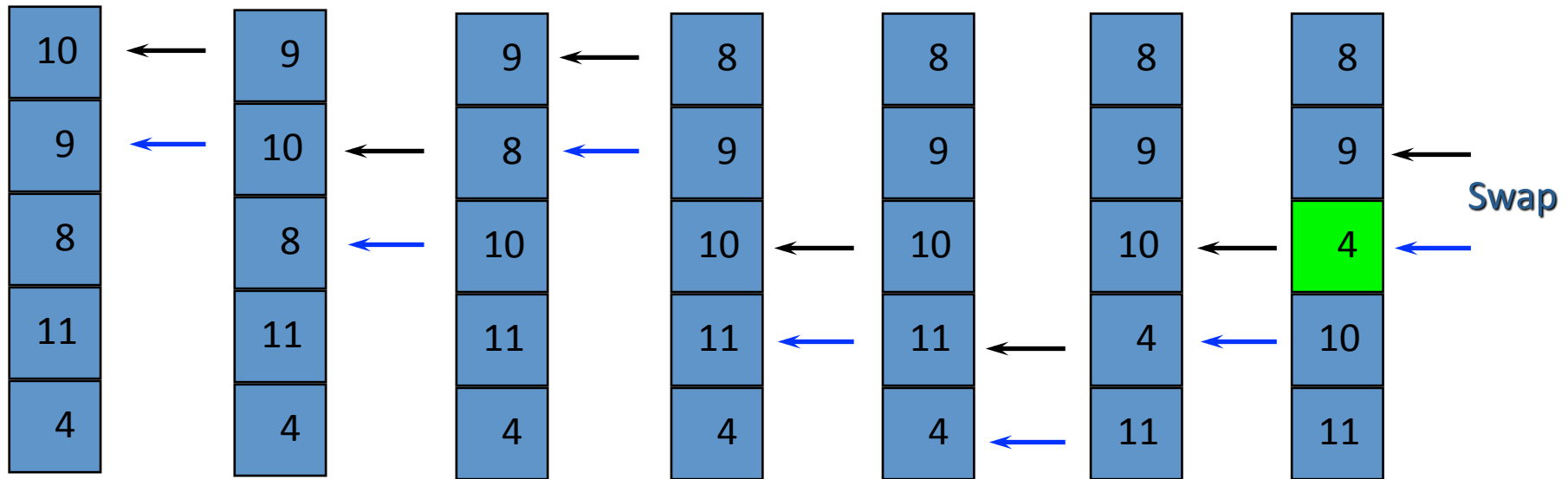
Bubble Sort



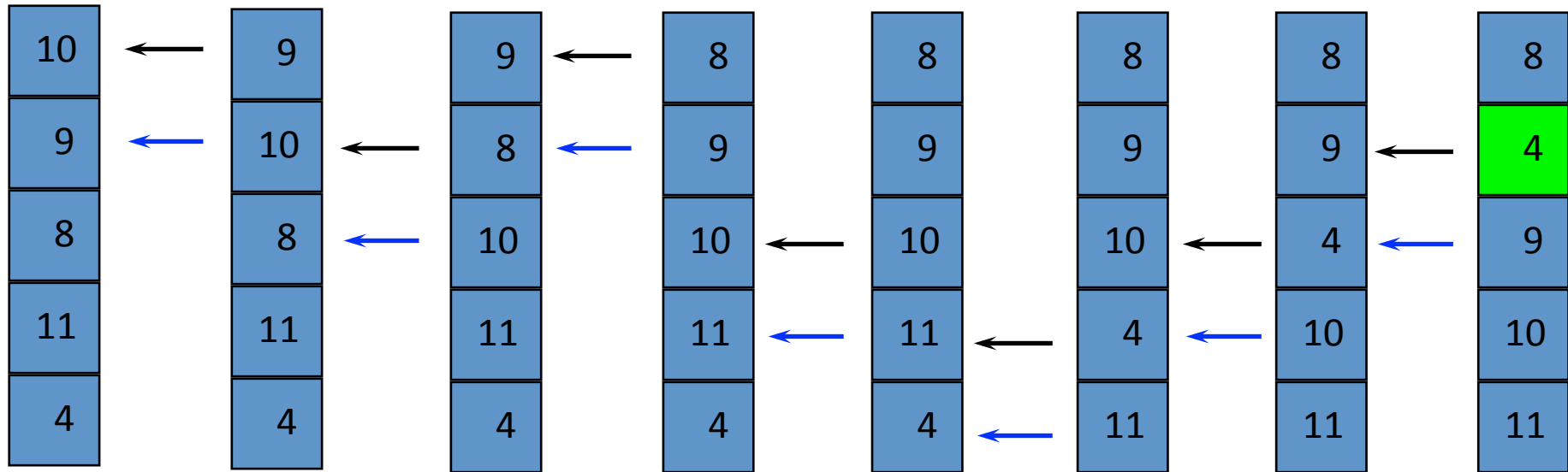
Bubble Sort



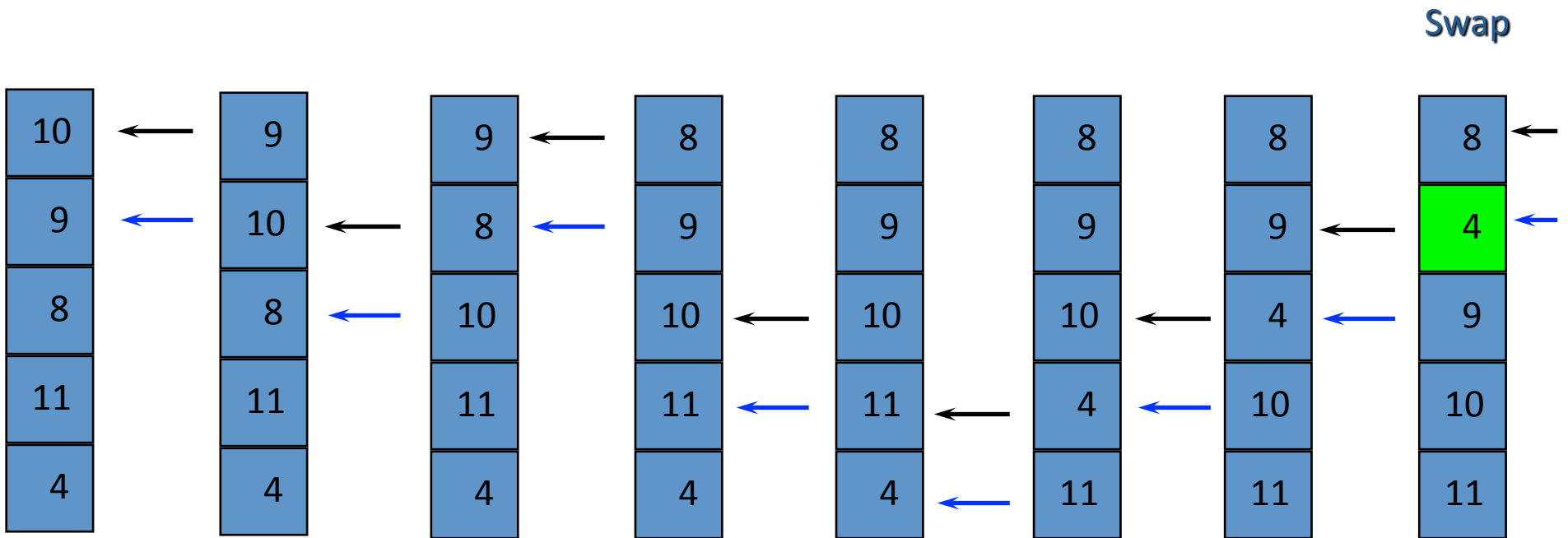
Bubble Sort



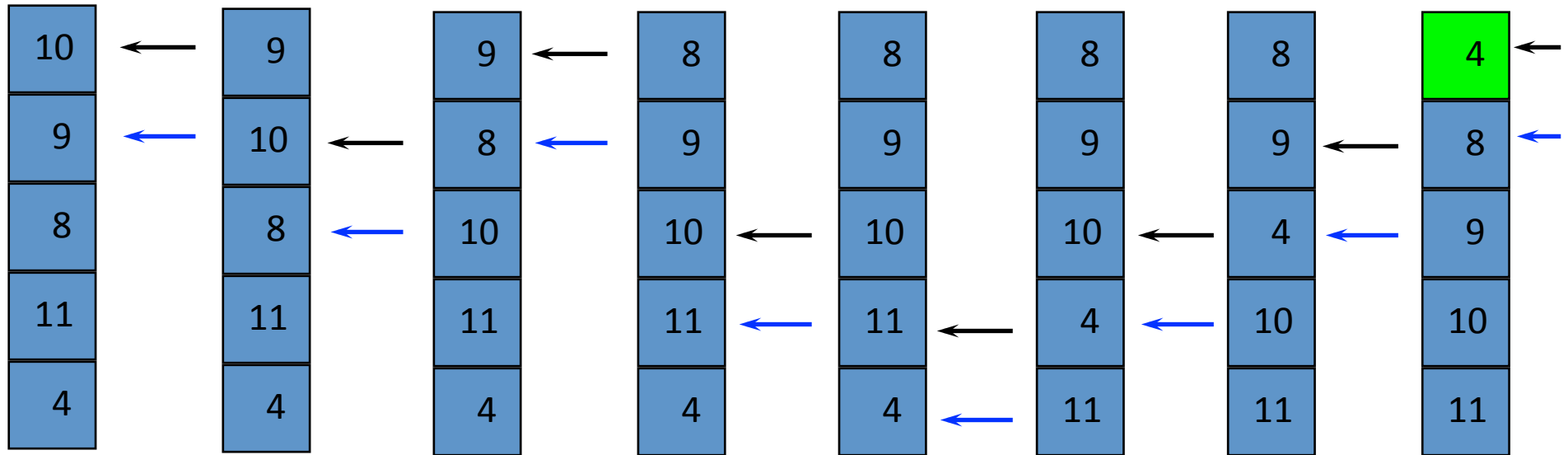
Bubble Sort



Bubble Sort



Bubble Sort



Implementation of Bubble_Sort()

```
int bubble_sort(int *a, int size) {  
    int i, j, temp;  
  
    for (i=0; i < size-1; i++) {  
        for (j=i; j >= 0; j--) {  
            if (a[j] > a[j+1]) {  
  
                /* swap */  
                temp = a[j+1];  
                a[j+1] = a[j];  
                a[j] = temp;  
            }  
        }  
    }  
}
```

Bubble Sort

- A few observations:
 - we don't usually sort numbers; we usually sort records with keys
 - the key can be a number
 - or the key could be a string
 - the record would be represented with a `struct`
 - The swap should be done with a function (so that a record can be swapped)
 - We can make the preceding algorithm more efficient. How? (hint: do we always have to bubble back to the top?)

Bubble Sort

- Exercise: implement these changes and write a driver program to test:
 - the original bubble sort
 - the more efficient bubble sort
 - the bubble sort with a swap function
 - the bubble sort with structures
 - compute the order of time complexity of the bubble sort

Selection Sort

- Example:

- Shaded elements are selected
- Boldface elements are in order

Initial Array

29	10	14	37	13
----	----	----	----	----

After 1st swap

29	10	14	13	37
----	----	----	----	-----------

After 2nd swap

13	10	14	29	37
----	----	----	-----------	-----------

After 3rd swap

13	10	14	29	37
----	----	-----------	-----------	-----------

After 4th swap

10	13	14	29	37
-----------	-----------	-----------	-----------	-----------

Selection Sort

- Assume we are sorting a list represented by an array A of n integer elements
- Selection sort algorithm in pseudo-code

```
last =  $n-1$ 
Do
    Select largest element from  $a[0..last]$ 
    Swap it with  $a[last]$ 
    last = last-1
While (last >= 1)
```

Selection Sort

```
typedef int DataType;

void selectionSort(DataType a[] , int n) {

    DataType temp;
    int index_of_largest, index, last;

    for(last= n-1; last >= 1; last--) {

        // select largest item in a[0..last]
        index_of_largest = 0;
        for(index=1; index <= last; index++) {
            if (a[index] > a[index_of_largest])
                index_of_largest = index;
        }

        // swap largest item with last element
        temp = a[index_of_largest];
        a[index_of_largest] = a[last]);
        a[last]) = temp;
    }
}
```

Quicksort

- The Quicksort algorithm was developed by C.A.R. Hoare. It has the best average behaviour in terms of complexity:

Average case: $O(n \log_2 n)$

Worst case: $O(n^2)$

Quicksort

- Given a list of elements
- take a partitioning element
- and create a (sub)list
 - such that all elements to the left of the partitioning element are less than it,
 - and all elements to the right of it are greater than it
- Now repeat this partitioning effort on each of these two sublists

Quicksort

- And so on in a recursive manner until all the sublists are empty, at which point the (total) list is sorted
- Partitioning can be effected simultaneously, scanning left to right and right to left, interchanging elements in the wrong parts of the list\
- The partitioning element is then placed between the resultant sublists (which are then partitioned in the same manner)

Implementation of Quicksort()

In pseudo-code first

If anything to be partitioned

choose a pivot

DO

scan from left to right until we find an element
> pivot: i points to it

scan from right to left until we find an element
< pivot: j points to it

IF $i < j$

exchange ith and jth element

WHILE $i \leq j$

Implementation of Quicksort()

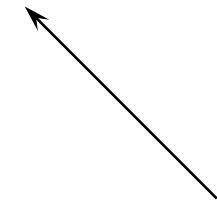
```
/* simple quicksort to sort an array of integers */  
  
void quicksort (int A[], int L, int R)  
{  
    int i, j, pivot;  
  
    /* assume A[R] contains a number > any element, */  
    /* i.e. it is a sentinel. */
```

Implementation of Quicksort()

```
if ( R > L) {
    i = L; j = R;
    pivot = A[i];
    do {
        while (A[i] <= pivot) i=i+1;
        while ((A[j] >= pivot) && (j>1)) j=j-1;
        if (i < j) {
            exchange(A[i],A[j]); /*between partitions*/
            i = i+1; j = j-1;
        }
    } while (i <= j);
    exchange(A[L], A[j]); /* reposition pivot */
    quicksort(A, L, j);
    quicksort(A, i, R);    /*includes sentinel*/
}
}
```

Quicksort

10	9	8	11	4	99
----	---	---	----	---	----



sentinel

Quicksort

10	9	8	11	4	99
----	---	---	----	---	----

QS (A, ,)

L:

R:

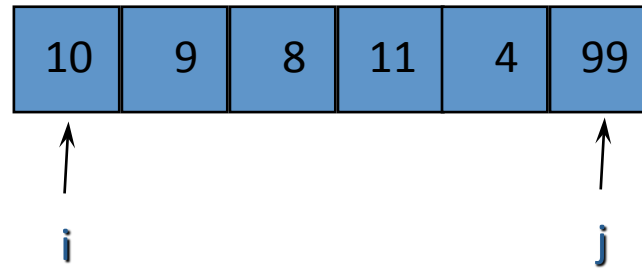
i:

j:

pivot:



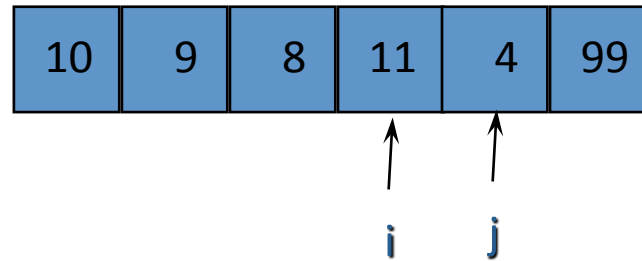
Quicksort



QS (A, 1, 6)

L: 1
R: 6
i: 1
j: 6
pivot: 10

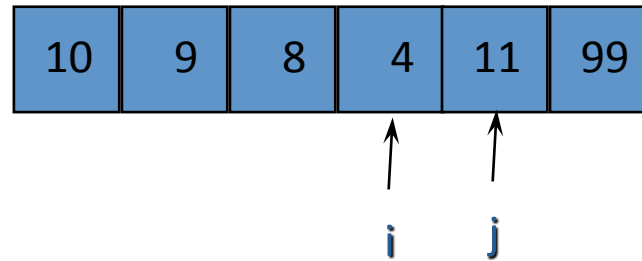
Quicksort



QS (A, 1, 6)

L: 1
R: 6
i: 1 2 3 4
j: 6 5
pivot: 10

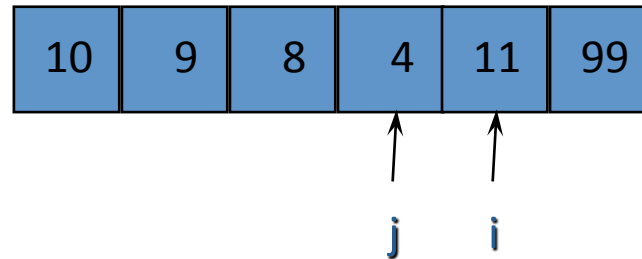
Quicksort



QS (A, 1, 6)

L: 1
R: 6
i: 1 2 3 4
j: 6 5
pivot: 10

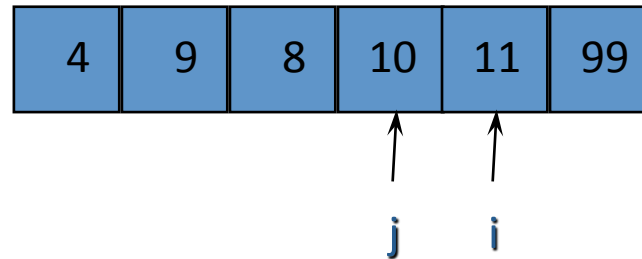
Quicksort



QS (A, 1, 6)

L: 1
R: 6
i: 1 2 3 4 5
j: 6 5 4
pivot: 10

Quicksort



QS (A, 1, 6)

L: 1
R: 6
i: 1 2 3 4 5
j: 6 5 4
pivot: 10

Quicksort

4	9	8	10	11	99
---	---	---	----	----	----

↑ ↑
i j

QS (A, 1, 6)

L: 1
R: 6
i: 1 2 3 4 5
j: 6 5 4
pivot: 10

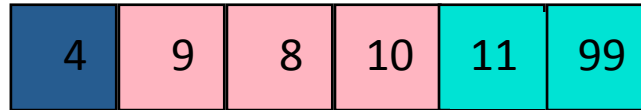
QS (A, 1, 4)

L: 1
R: 4
i:
j:
pivot: 4

QS (A, 5, 6)

L: 5
R: 6
i:
j:
pivot: 11

Quicksort



↑
i

↑
j

QS (A, 1, 1)

L: 1
R: 1
i:
j:
pivot: 4

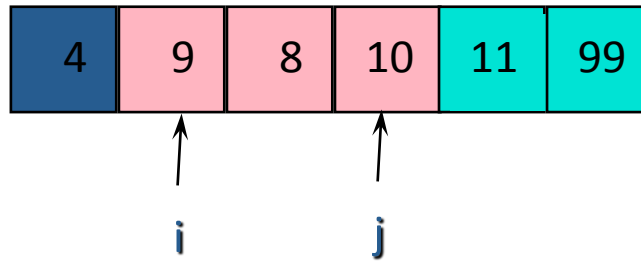
QS (A, 2, 4)

L: 2
R: 4
i:
j:
pivot: 9

QS (A, 5, 6)

L: 5
R: 6
i: 5
j: 6
pivot: 11

Quicksort



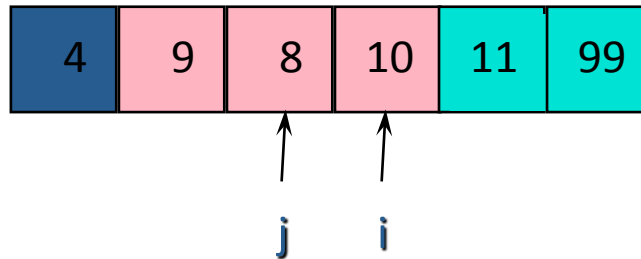
QS (A, 2, 4)

L: 2
R: 4
i: 2
j: 4
pivot: 9

QS (A, 5, 6)

L: 5
R: 6
i: 5
j: 6
pivot: 11

Quicksort



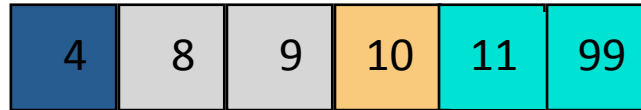
QS (A, 2, 4)

L: 2
R: 4
i: 2 3 4
j: 4 3
pivot: 9

QS (A, 5, 6)

L: 5
R: 6
i: 5
j: 6
pivot: 11

Quicksort



↑ ↑
i j

QS (A, 2, 4)

L: 2
R: 4
i: 2 3 4
j: 4 3
pivot: 9

QS (A, 2, 3)

L: 2
R: 3
i:
j:
pivot: 8

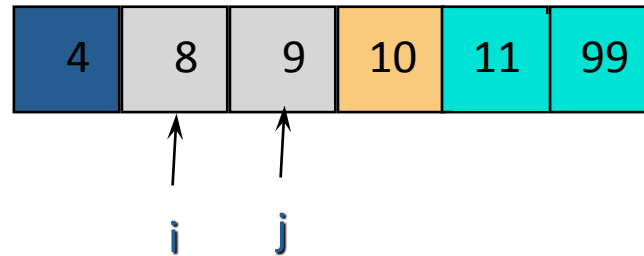
QS (A, 4, 4)

L: 4
R: 4
i:
j:
pivot: 10

QS (A, 5, 6)

L: 5
R: 6
i: 5
j: 6
pivot: 11

Quicksort



QS (A, 2, 3)

L: 2
R: 3
i: 2
j: 3
pivot: 8

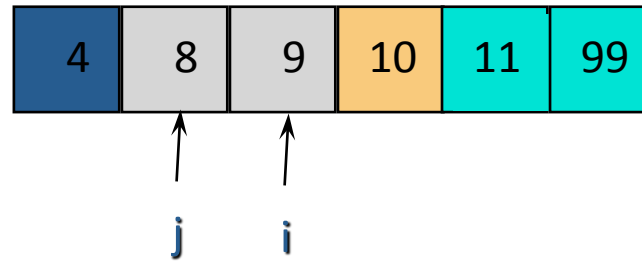
QS (A, 4, 4)

L: 4
R: 4
i:
j:
pivot: 10

QS (A, 5, 6)

L: 5
R: 6
i: 5
j: 6
pivot: 11

Quicksort



QS (A, 2, 3)

L: 2
R: 3
i: 2 3
j: 3 2
pivot: 8

QS (A, 4, 4)

L: 4
R: 4
i:
j:
pivot: 10

QS (A, 5, 6)

L: 5
R: 6
i: 5
j: 6
pivot: 11

Quicksort



↑ ↑
i j

QS (A, 2, 3)

QS (A, 2, 2)

QS (A, 3, 3)

QS (A, 4, 4)

QS (A, 5, 6)

L: 2
R: 3
i: 2 3
j: 3 2
pivot: 8

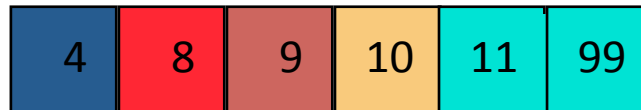
L: 2
R: 2
i:
j:
pivot: 8

L: 3
R: 3
i:
j:
pivot: 9

L: 4
R: 4
i:
j:
pivot: 10

L: 5
R: 6
i: 5
j: 6
pivot: 11

Quicksort



↑ ↑
i j

QS (A, 4, 4)

L: 4
R: 4
i:
j:
pivot: 10

QS (A, 5, 6)

L: 5
R: 6
i: 5
j: 6
pivot: 11

Quicksort



↑ ↑
i j

QS (A, 5, 6)

L: 5
R: 6
i: 5
j: 6
pivot: 11

QS (A, 5, 5)

L: 5
R: 5
i:
j:
pivot: 11

QS (A, 6, 6)

L: 6
R: 6
i:
j:
pivot: 99