

# Data Structures and Algorithms for Engineers

## Module 5: Lists

Lecture 4: Queues. Implementation using List ADT. Comparison of order of complexity. Dedicated ADT. Circular queues. Queue applications.

David Vernon  
Carnegie Mellon University Africa

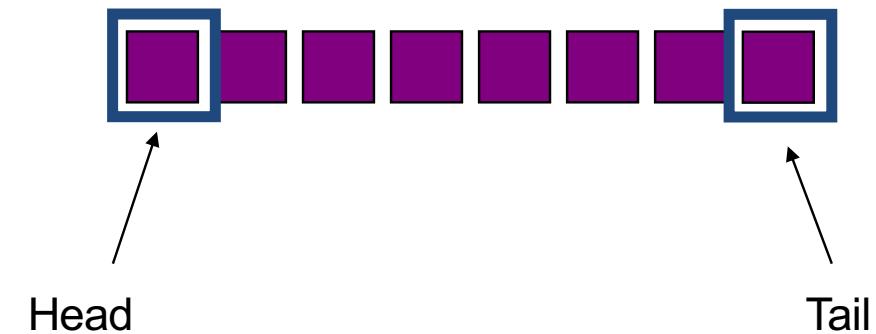
[vernon@cmu.edu](mailto:vernon@cmu.edu)  
[www.vernon.eu](http://www.vernon.eu)

# Queues

A queue is another special type of list

- Insertions are made at one end, called the tail of the queue
- Deletions take place at the other end, called the head
- Thus, the last one added is always the last one available for deletion
- Also referred to as a FiFO list (First In First Out)

# Queues



# Queue Operations

Declare: → Q :

The function value of **Declare(Q)** is an empty queue

# Queue Operations

Empty:  $\rightarrow Q :$

The function **Empty** causes the queue to be emptied and it returns position **End(Q)**



# Queue Operations

IsEmpty:  $Q \rightarrow B$  :

The function value **IsEmpty(Q)** is true if Q is empty;  
otherwise it is false

# Queue Operations

Head:  $Q \rightarrow E$  :

The function value  $\text{Head}(Q)$  is the first element in the list;

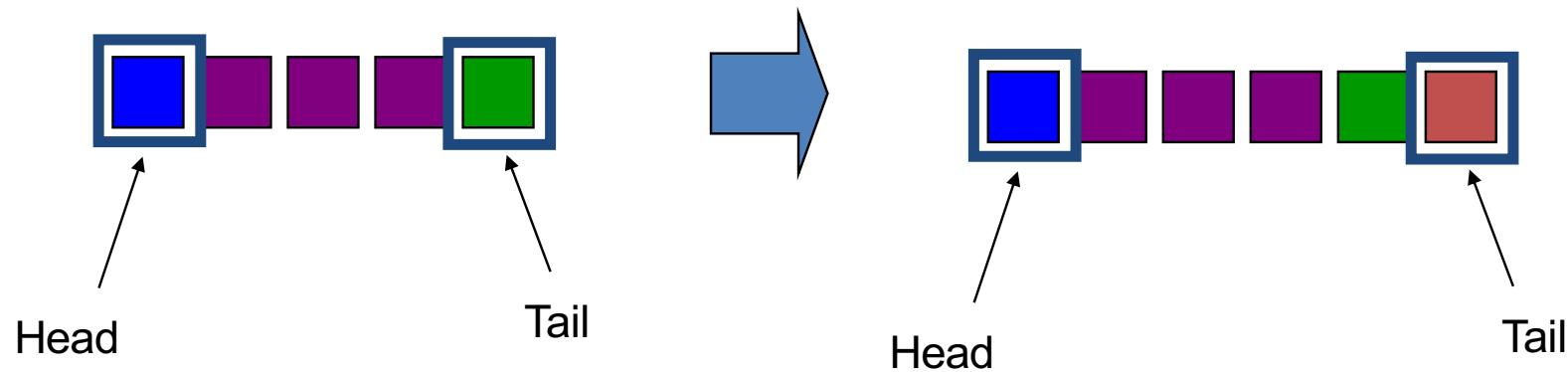
if the queue is empty, the value is undefined

# Queue Operations

Enqueue:  $E \times Q \rightarrow Q$  :

**Enqueue( $e, Q$ )**

Add an element  $e$  to the tail of the queue

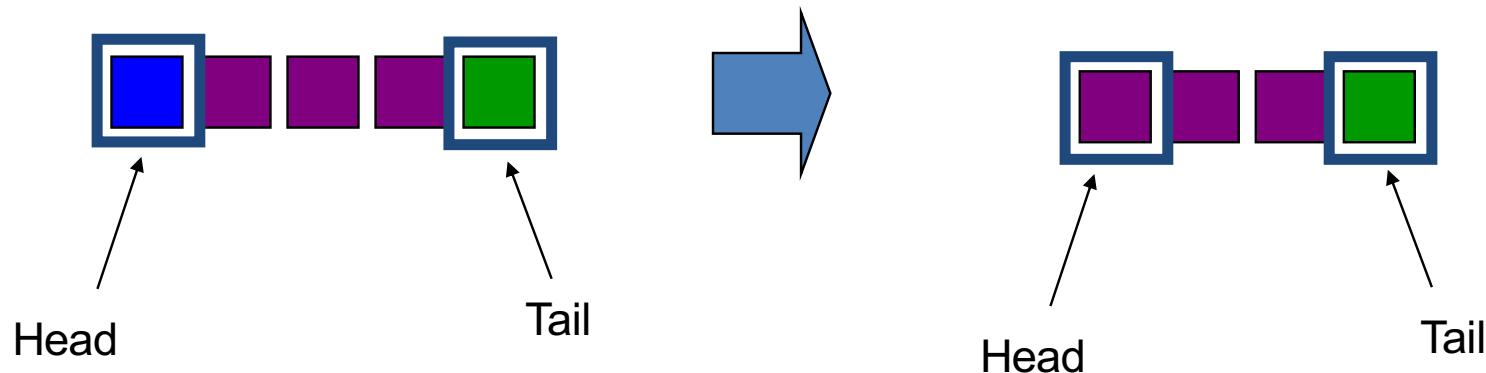


# Queue Operations

Dequeue:  $Q \rightarrow E$  :

## Dequeue(Q)

Remove the element from the head of the queue: i.e., return the first element and delete it from the queue



# Queue Operations

- All these operations can be directly implemented using the LIST ADT operations on a queue Q
- Again, it may be more efficient to use a dedicated implementation
- And, again, it depends what you want: code efficiency or software re-use (i.e, utilization efficiency)

# Queue Operations

Declare(Q)

Empty(Q)

Head(Q)

    Retrieve(First(Q), Q)

Enqueue(e, Q)

    Insert(e, End(Q), Q)

Dequeue(Q)

    Retrieve(First(Q), Q)

    Delete(First(Q), Q)

# Queue Errors

- Queue **overflow** errors occur when you attempt to **enqueue()** an element in a queue that is **full**
- Queue **underflow** errors occur when you attempt to **dequeue()** an element from an empty queue
- Your ADT implementation should provide guards that catch these errors

# Queue Implementation

- The List ADT can be implemented
  - As an array
  - As a linked-list
- So, therefore, so can the Queue ADT
- What are the advantages and disadvantages of these two options?
- When would you pick one implementation over the other?

# Queue Operations

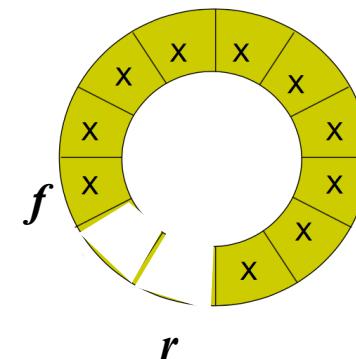
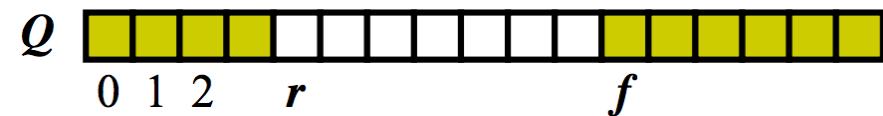
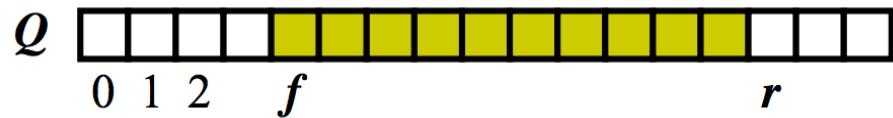
	Array	Linked-List
Declare(Q)	O(1)	O(1)
Empty(Q)	O(1)	O( $n$ )
Head(Q)	O(1)	O(1)
Retrieve(First(Q), Q)		
Enqueue(e, Q)	O(1)	O( $n$ ) ... why?
Insert(e, End(Q), Q)		
Dequeue(Q)	O( $n$ ) ... why?	O(1)
Retrieve(First(Q), Q)		
Delete(First(Q), Q)		

# Queue Implementation

- Reusing the List ADT involves some compromises
- Alternative is to create a new Queue ADT
  - With an implementation that avoids these compromises

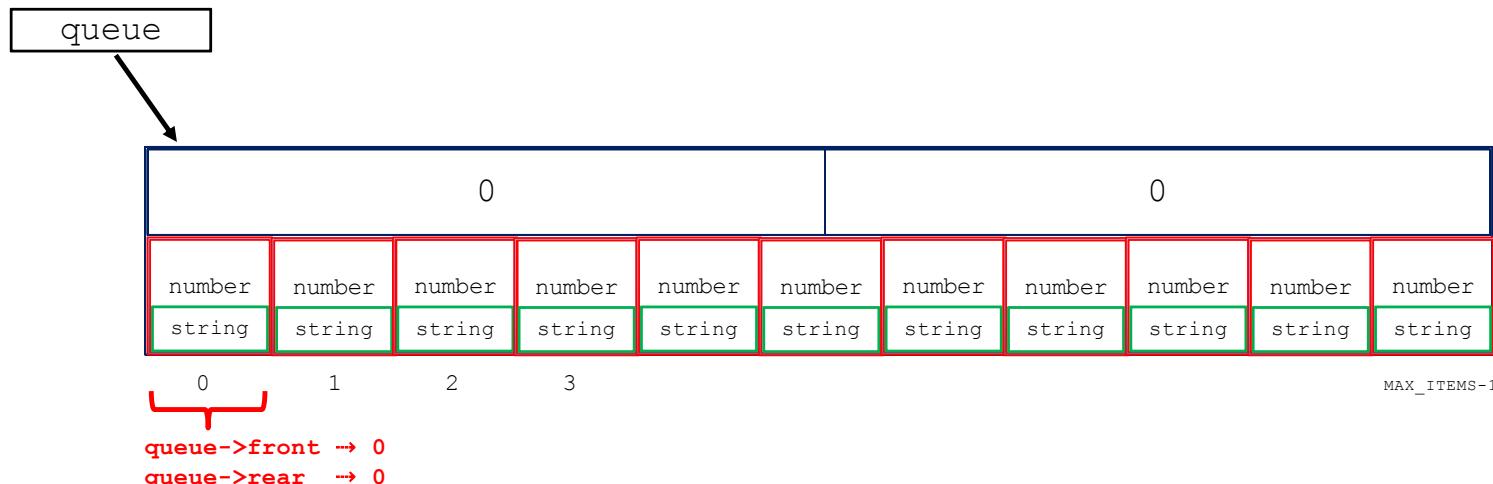
# Queue Implementation

```
typedef struct {
    int front;
    int rear;
    ITEM_TYPE items[MAX_ITEMS];
} QUEUE_TYPE;
```



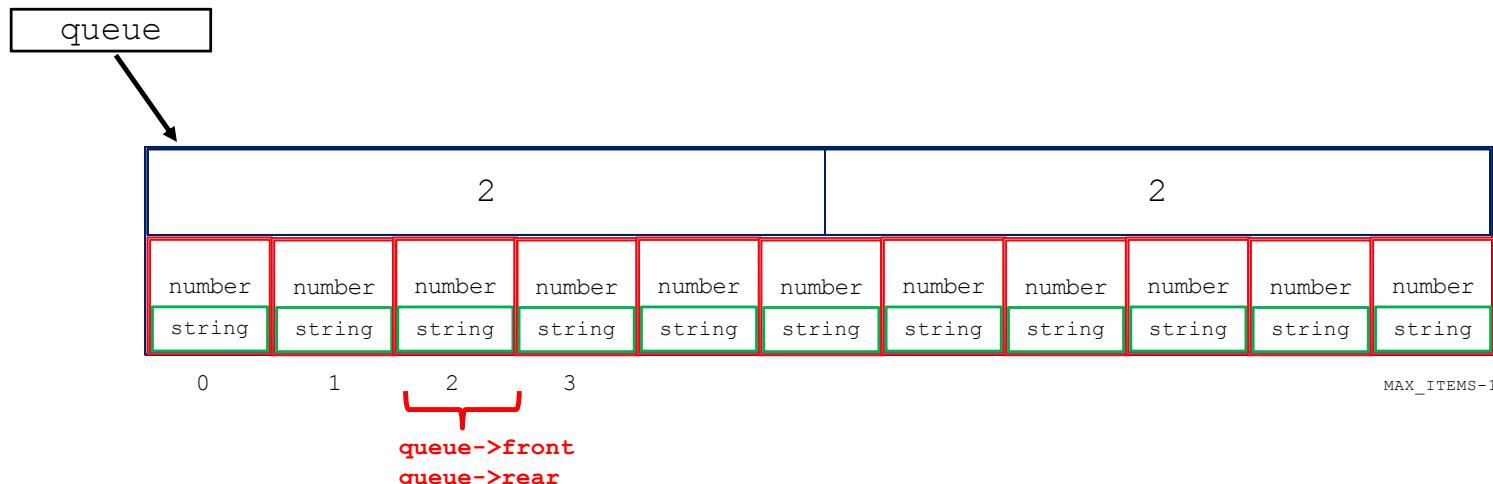
# Queue Implementation

```
void empty(QUEUE_TYPE *queue) {  
    queue->front = 0;  
    queue->rear = 0;  
    return (end(queue));  
}
```



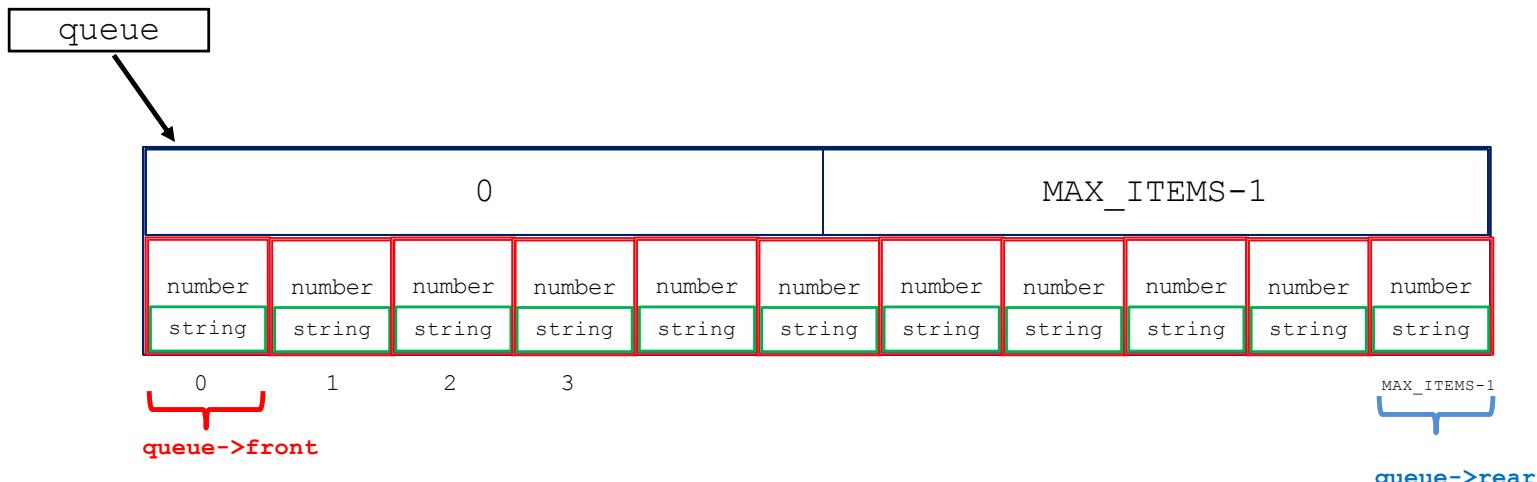
# Queue Implementation

```
bool is_empty(QUEUE_TYPE *queue) {  
    if (queue->front == queue->rear)  
        return(true);  
    else  
        return(false)  
}
```



# Queue Implementation

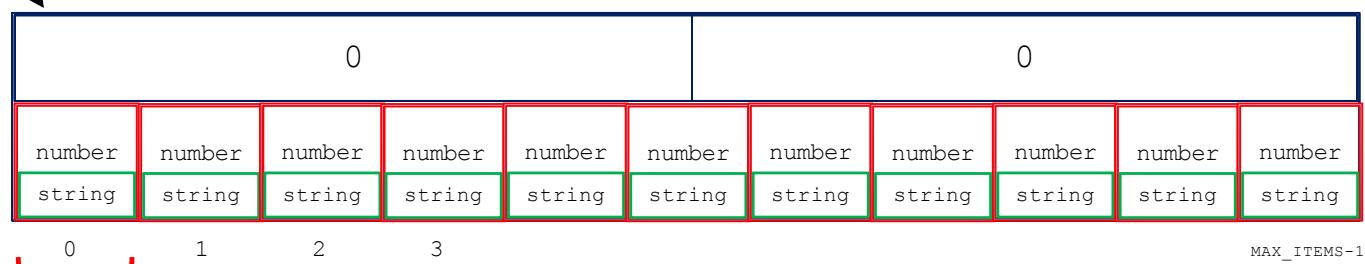
```
int is_full(QUEUE_TYPE *queue) {  
    if ((queue->rear + 1) % MAX_ITEMS == queue->front)  
        return(TRUE);  
    else  
        return(FALSE);  
}
```



# Queue Implementation

```
void enqueue(ITEM_TYPE e, QUEUE_TYPE *queue) {  
    if (!is_full(queue)) {  
        queue->items[queue->rear] = e;  
        queue->rear = (queue->rear + 1) % MAX_ITEMS;  
    }  
    else {  
        error("Queue overflow: queue is already full");  
    }  
}
```

queue

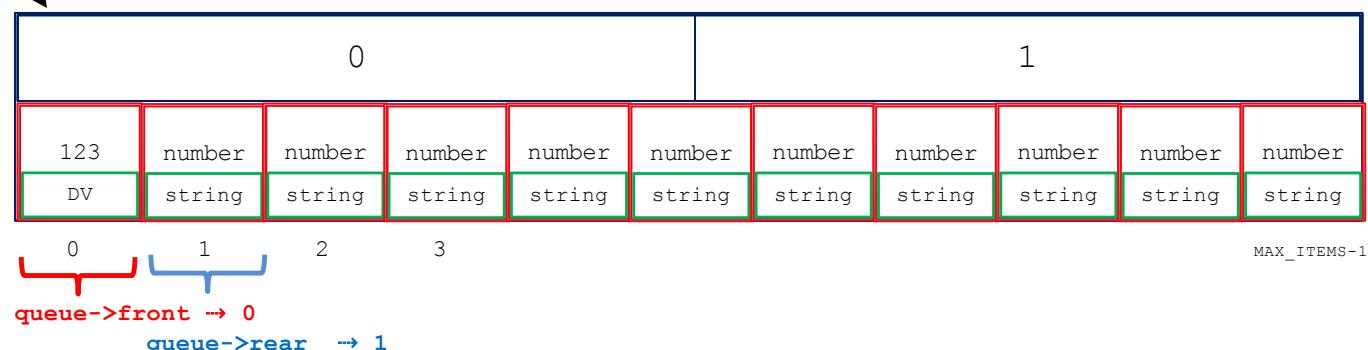


`queue->front → 0`  
`queue->rear → 0`

# Queue Implementation

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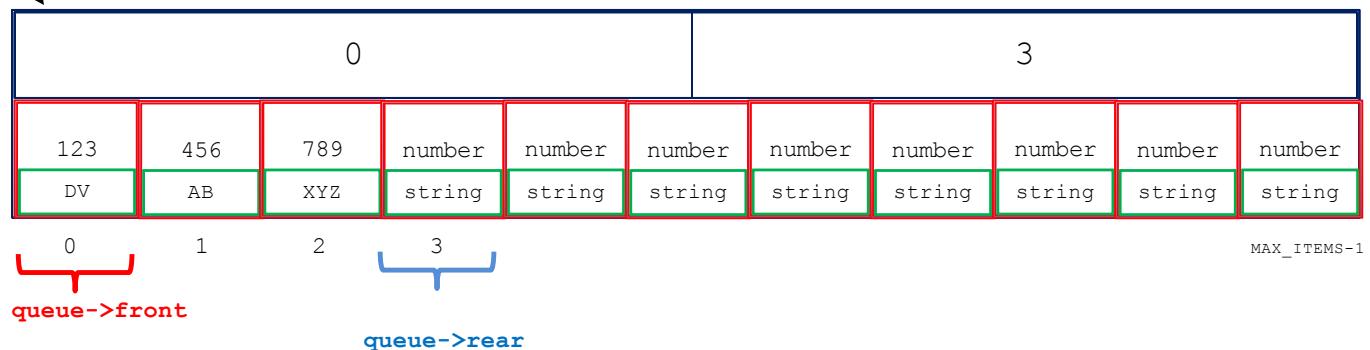
queue



# Queue Implementation

```
void dequeue(ITEM_TYPE *e, QUEUE_TYPE *queue) {
    if (!is_empty(queue)) {
        *e = queue->items[queue->front];
        queue->front = (queue->front+1) % MAX_ITEMS;
    }
    else {
        error("Queue underflow: queue is empty");
    }
}
```

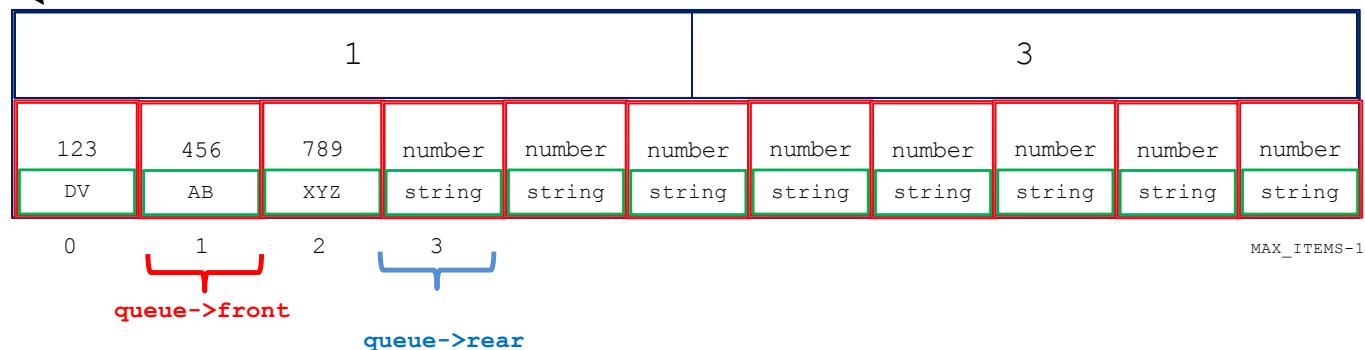
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queue



# Queue Implementation

- Can you see a particular problem with the linked-list implementation?
- How would you fix it?

# Queue Operations

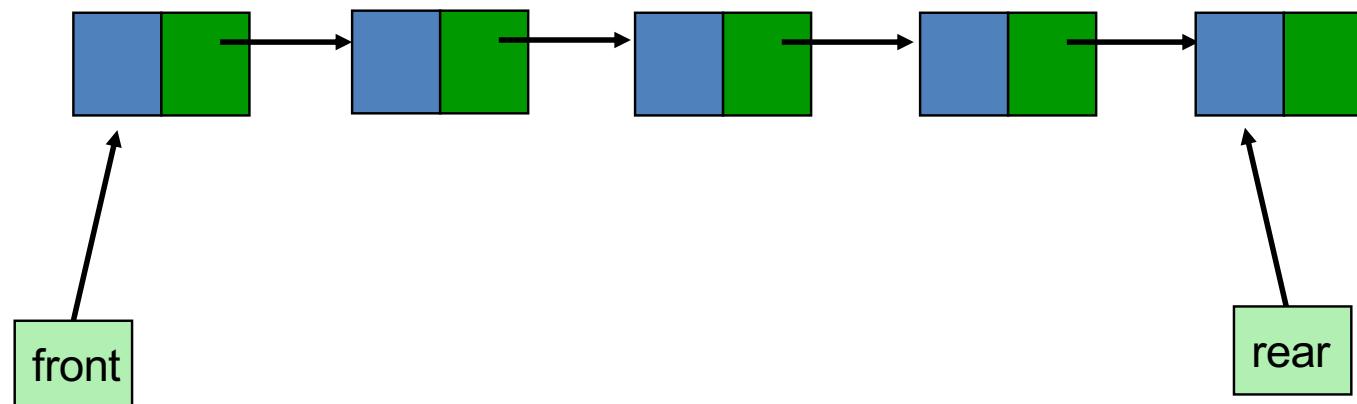
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Delete(First(Q), Q)		

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# Queue Applications

- Scheduling/waiting for system service queues
- Resource queues – provide coordinated access to shared resources
- Message queues (Buffer)
- Multi-queues and priority queues