Data Structures and Algorithms for Engineers

Module 6: Trees

Lecture 2: Binary Search Tree

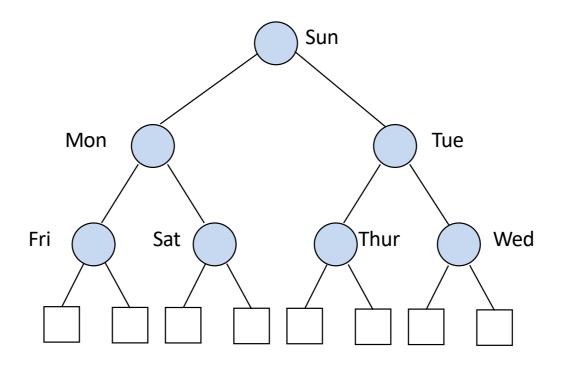
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- A Binary Search Tree (BST) is a special type of binary tree
 - it represents information is an ordered format
 - A binary tree is binary search tree if for every node w,
 all keys in the left subtree of w have values less than the key of w
 all keys in the right subtree have values greater than key of w.

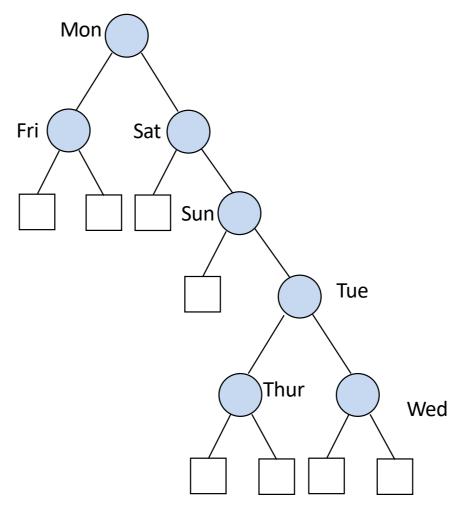
Definition: A binary search tree T is a binary tree; either it is empty or each node in the tree contains an identifier and:

- all keys in the left subtree of T are less (numerically or alphabetically) than the identifier in the root node T
- all identifiers in the right subtree of T are greater than the identifier in the root node T
- The left and right subtrees of T are also binary search trees.



- The main point to notice about such a tree is that, if traversed inorder, the keys of the tree (i.e., its data elements) will be encountered in a sorted fashion
- Furthermore, efficient searching is possible using the binary search technique
 - search time is $O(\log_2 n)$

It should be noted that several binary search trees are possible for a given data set, e.g., consider the following tree:



Binary Search Trees Fri Mon Sat Sun Thur Tue Wed

Let us consider how such a situation might arise

Construct a binary search tree:

- Assume we are building a binary search tree of words
- Initially, the tree is null, i.e. there are no nodes in the tree
- The first word is inserted as a node in the tree as the root, with no children

On insertion of the second word, we check to see if it is the same as the key in the root, less than it, or greater than it

- If it is the same, no further action is required (duplicates are not allowed)
- If it is less than the key in the current node, move to the left subtree and compare again
- If the left subtree does not exist, then the word does not exist, and it is inserted as a new node on the left

- If, on the other hand, the word was greater than the key in the current node, move to the right subtree and compare again
- If the right subtree does not exist, then the word does not exist, and it is inserted as a new node on the right
- This insertion can most easily be effected in a recursive manner

- The point here is that the structure of the tree depends on the order in which the data is inserted in the list
- If the words are entered in sorted order, then the tree will degenerate to a
 1-D list

BST Operations

Insert: $E \times BST \rightarrow BST$:

The function value Insert(e,T) is the BST T with the element e inserted as a leaf node; if the element already exists, no action is taken

NO WINDOW!!!

BST Operations

Delete: $E \times BST \rightarrow BST$:

The function value Delete(e, T) is the BST T with the element e deleted; if the element is not in the BST exists, no action is taken.

NO WINDOW!!!

Implementation of Insert(e, T)

- If T is empty (i.e., T is NULL)
 - create a new node for e
 - make T point to it
- If T is not empty
 - if e < element at root of T</p>
 - Insert e in left child of T: Insert(e, T(1))
 - if e > element at root of T
 - Insert e in right child of T: Insert(e, T(2))

First, we must locate the element e to be deleted in the tree

- if e is at a leaf node
 - we can delete that node and be done
- if e is at an interior node at w
 - we can't simply delete the node at w as that would disconnect its children
- if the node at w has only one child
 - we can replace that node with its child

- if the node at w has two children
 - we replace the node at w with the lowest-valued element among the descendents of its right child
 - this is the left-most node of the right sub-tree
 - It is useful to have a function DeleteMin() which removes the smallest element from a non-empty tree and

returns the value of the element removed

- If T is not empty
 - if e < element at root of T

Delete e from left child of T: Delete(e, T(1))

- if e > element at root of T

Delete e from right child of T: Delete(e, T(2))

— if e = element at root of T and both children are empty

Remove T

if e = element at root of T and left child is empty

Replace T with T(2)

if e = element at root of T and right child is empty

Replace T with T(1)

if e = element at root of T and neither child is empty

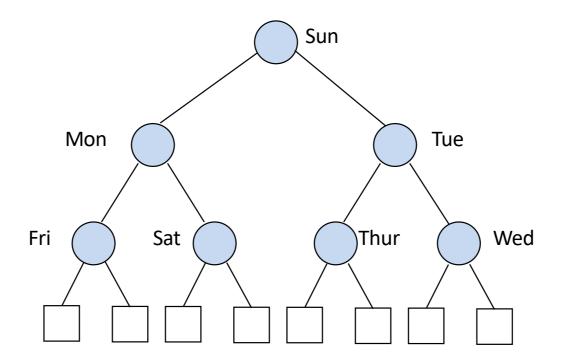
Replace T with left-most node of $T(2) \leftarrow "left-most node in right sub-tree!"$

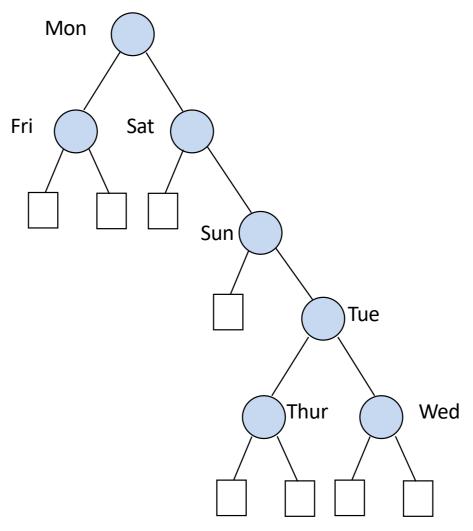
What if the left-most node in the right sub-tree has two (interior node) children?

lt can't!

If it did, it wouldn't be the left-most node ...

because there would be a node on it's left!



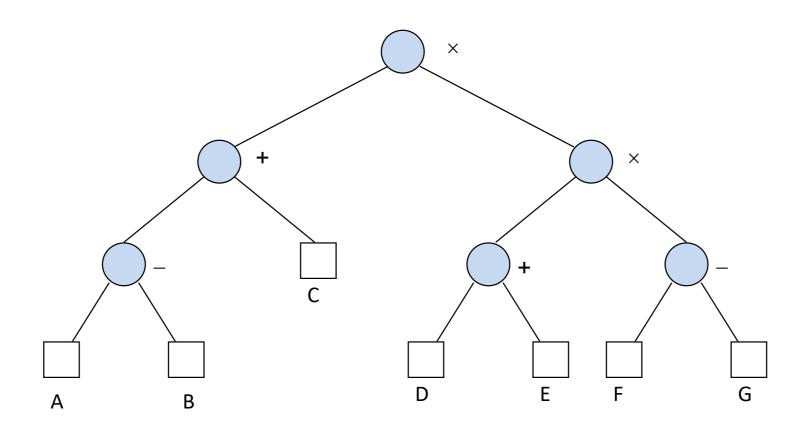


Tree Traversals

- To perform a traversal of a data structure, we use a method of visiting every node in some predetermined order
- Traversals can be used
 - to test data structures for equality
 - to display a data structure
 - to construct a data structure of a given size
 - to copy a data structure

- There are 3 depth-first traversals
 - Inorder
 - Postorder
 - Preorder
- For example, consider the expression tree:

Example: Expression Tree



Inorder traversal

$$A-B+CxD+ExF-G$$

Postorder traversal

$$AB-C+DE+FG-xx$$

Preorder traversal

$$x + -ABCx + DE - FG$$

• The parenthesised Inorder traversal

$$((A - B) + C) \times ((D + E) \times (F - G))$$

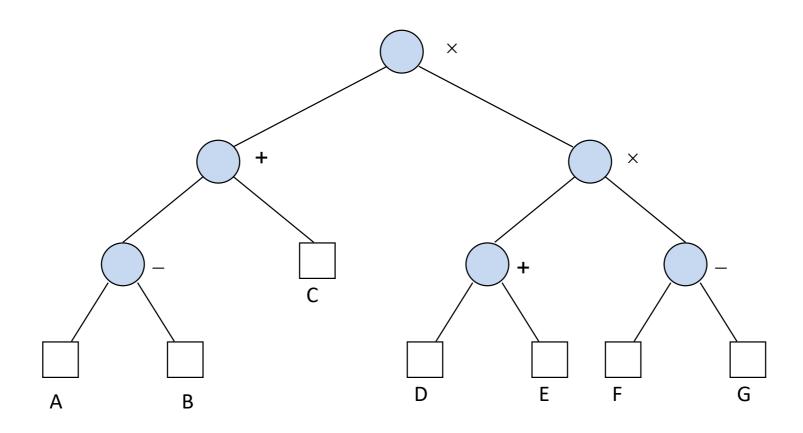
This is the infix expression corresponding to the expression tree

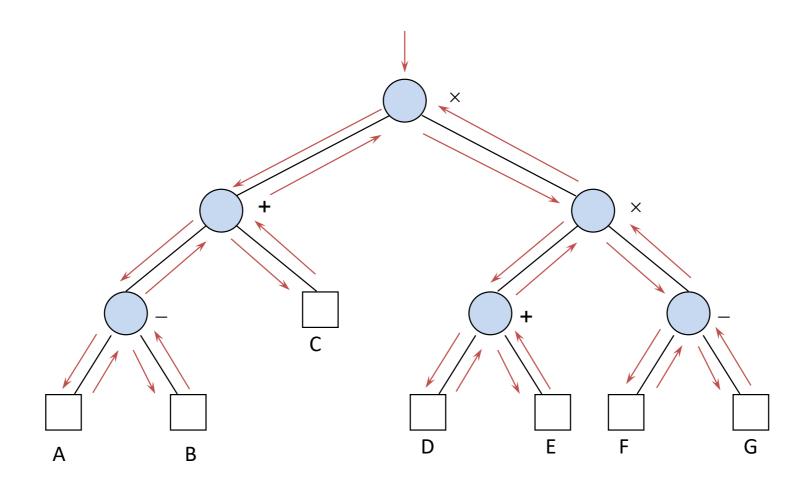
- Postorder traversal gives a postfix expression
- Preorder traversal gives a prefix expression

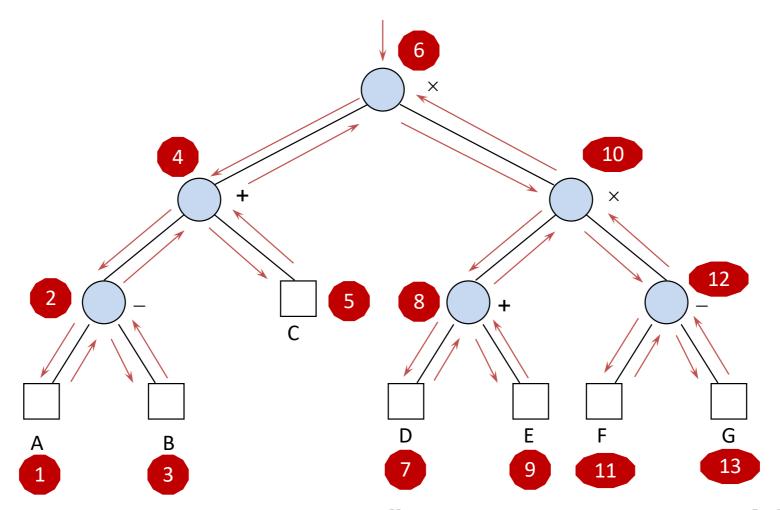
Recursive definition of inorder traversal

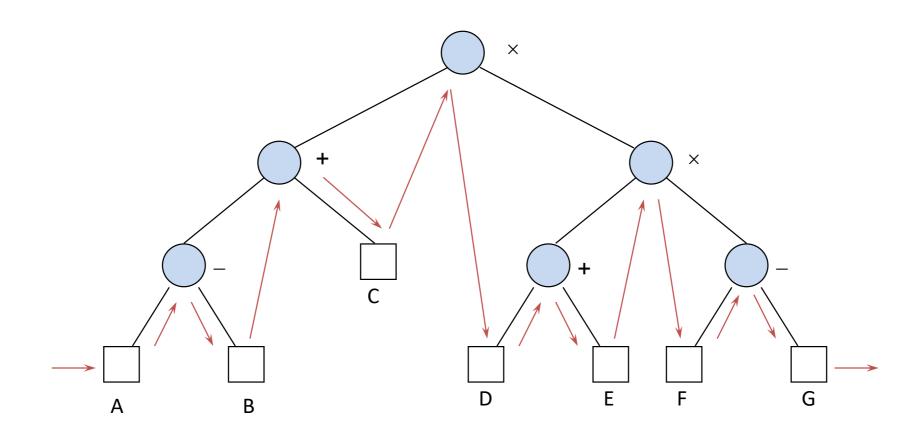
Given a binary tree T

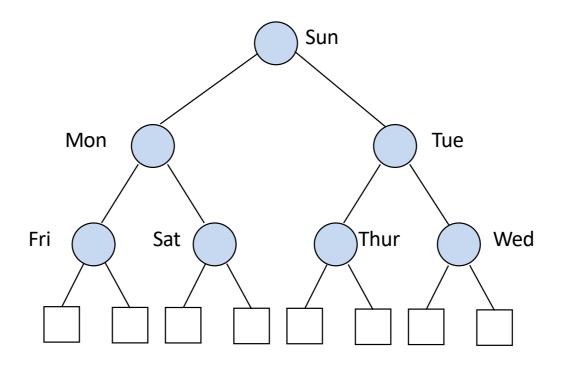
```
visit the external node
otherwise
perform an inorder traversal of Left(T)
visit the root of T
perform an inorder traversal of Right(T)
```

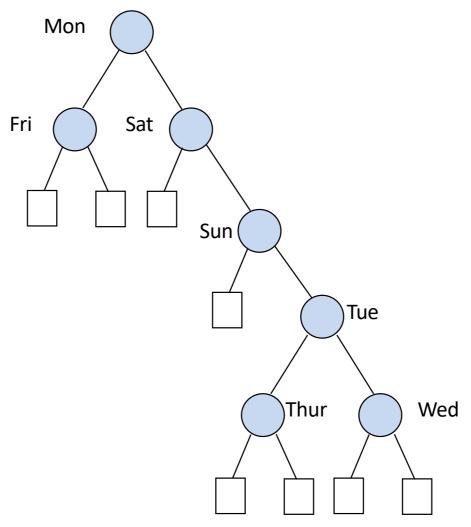












Recursive definition of postorder traversal

```
Given a binary tree T

if T is empty

visit the external node

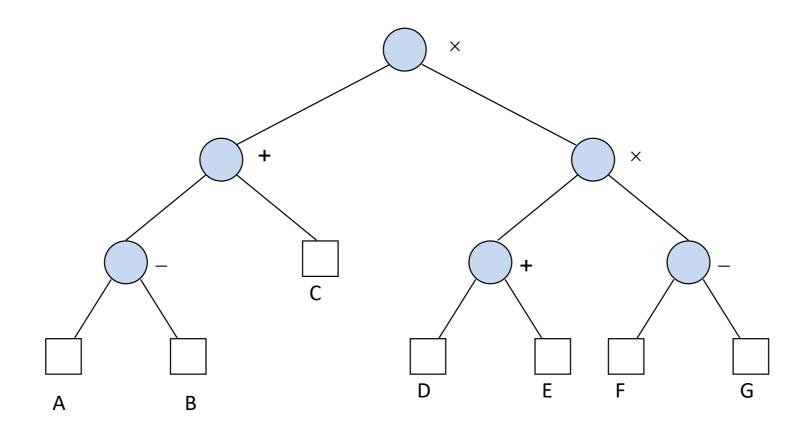
otherwise

perform a postorder traversal of Left(T)

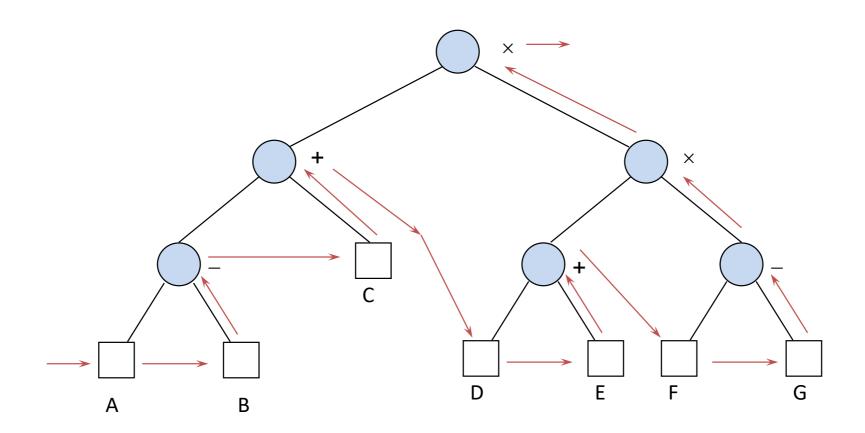
perform a postorder traversal of Right(T)

visit the root of T
```

Example: Postorder Traversal



Example: Postorder Traversal



Depth-First Traversals

Recursive definition of preorder traversal

```
Given a binary tree T

if T is empty

visit the external node

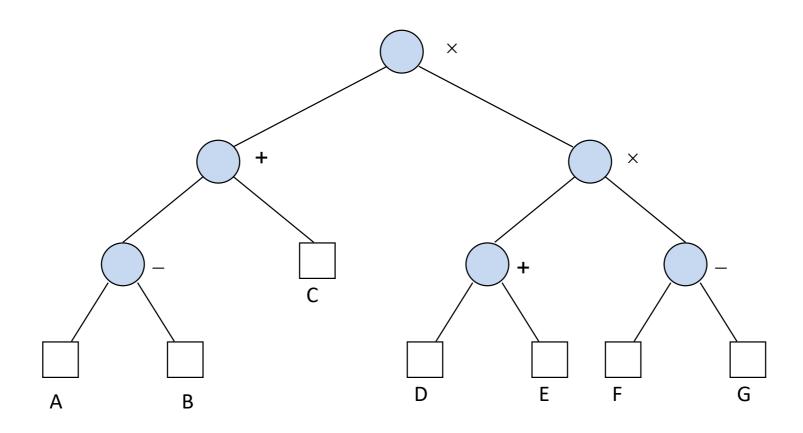
otherwise

visit the root of T

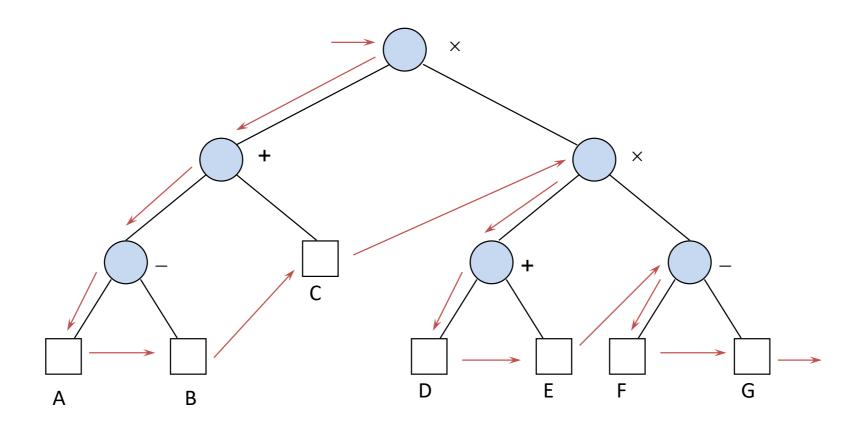
perform a preorder traversal of Left(T)

perform a preorder traversal of Right(T)
```

Example: Preorder Traversal



Example: Preorder Traversal



BST Implementation

```
typedef struct {
        int number;
        char *string;
    } ELEMENT_TYPE;

typedef struct node *NODE_TYPE;

typedef struct node {
        ELEMENT_TYPE element;
        NODE_TYPE left, right;
    } NODE;

typedef NODE_TYPE BINARY_TREE_TYPE;

typedef BINARY_TREE_TYPE WINDOW_TYPE;
```

```
int main() {
   ELEMENT_TYPE e;
   BINARY TREE TYPE tree;
   initialize(&tree);
   print(tree);
   assign_element_values(&e, 3, "...");
   insert(e, &tree);
   print(tree);
   assign_element_values(&e, 1, "+++");
   insert(e, &tree);
   print(tree);
   assign_element_values(&e, 5, "---");
   insert(e, &tree);
   print(tree);
   assign_element_values(&e, 2, ";;;");
   insert(e, &tree);
   print(tree);
   assign element values(&e, 4, "***");
   insert(e, &tree);
   print(tree);
   assign element values(&e, 6, "000");
   insert(e, &tree);
   print(tree);
   assign_element_values(&e, 3, "...");
   delete element(e, &tree);
   print(tree);
```

```
/*** initialize a tree ***/
void initialize(BINARY_TREE_TYPE *tree) {
   static bool first_call = true;
  /* we don't know what value *tree has when the program is launched
   /* so we have to be careful not to dereference it
                                                                            */
   /* if it's the first call to initialize, there is no tree to be deleted */
                                                                            */
   /* and we just set *tree to NULL
   if (first call) {
      first_call = false;
      *tree = NULL;
   else {
      if (*tree != NULL) postorder_delete_nodes(*tree);
      *tree = NULL;
```

```
/*** insert an element in a tree ***/
BINARY TREE TYPE *insert(ELEMENT TYPE e, BINARY TREE TYPE *tree ) {
  WINDOW_TYPE temp;
  if (*tree == NULL) {
     /* we are at an external node: create a new node and insert it */
     if ((temp = (NODE TYPE) malloc(sizeof(NODE))) == NULL)
         error("function insert: unable to allocate memory");
     else {
        temp->element = e;
        temp->left = NULL;
        temp->right = NULL;
         *tree = temp;
  else if (e.number < (*tree)->element.number) { /* assume the number field is the key */
     insert(e, &((*tree)->left));
  else if (e.number > (*tree)->element.number) {
     insert(e, &((*tree)->right));
  /* if e.number == (*tree)->element.number, e already is in the tree so do nothing */
  return(tree);
```

```
/*** returns & deletes the smallest node in a tree (i.e. the left-most node) */
ELEMENT TYPE delete min(BINARY TREE TYPE *tree) {
  ELEMENT_TYPE e;
  BINARY TREE TYPE p;
  if ((*tree)->left == NULL) {
      /* tree points to the smallest element */
      e = (*tree)->element;
      /* replace the node pointed to by tree by its right child */
      p = *tree;
      *tree = (*tree)->right;
     free(p);
      return(e);
  else {
      /* the node pointed to by tree has a left child */
      return(delete_min(&((*tree)->left)));
```

```
/*** delete an element in a tree ***/
BINARY_TREE_TYPE *delete_element(ELEMENT_TYPE e, BINARY_TREE_TYPE *tree) {
  BINARY TREE TYPE p;
  if (*tree != NULL) {
     if (e.number < (*tree)->element.number) /* assume element.number is the */
        delete element(e, &((*tree)->left)); /* key
     else if (e.number > (*tree)->element.number)
        delete element(e, &((*tree)->right));
     else if (((*tree)->left == NULL) && ((*tree)->right == NULL)) {
        /* leaf node containing e - delete it */
        p = *tree;
        free(p);
        *tree = NULL;
```

```
else if ((*tree)->left == NULL) {
      /* internal node containing e and it has only a right child */
      /* delete it and make treepoint to the right child
      p = *tree;
      *tree = (*tree)->right;
      free(p);
   else if ((*tree)->right == NULL) {
      /* internal node containing e and it has only a left child */
      /* delete it and make treepoint to the left child
      p = *tree;
      *tree = (*tree)->left;
      free(p);
   else {
      /* internal node containing e and it has both left and right child */
      /* replace it with leftmost node of right sub-tree
      (*tree)->element = delete min(&((*tree)->right));
return(tree);
```

```
/*** inorder traversal of a tree, printing node elements **/
int inorder(BINARY_TREE_TYPE tree, int n) {
  int i;
  if (tree != NULL) {
      inorder(tree->left, n+1);
      for (i=0; i<n; i++) printf(" ");</pre>
      printf("%d %s\n", tree->element.number, tree->element.string);
      inorder(tree->right, n+1);
  return(0);
```

```
postorder
/*** inorder traversal of a tree, deleting node elements **/
int postorder_delete_nodes(BINARY_TREE_TYPE tree) {
   if (tree != NULL) {
      postorder_delete_nodes(tree->left);
      postorder_delete_nodes(tree->right);
      free(tree);
   }
   return(0);
}
```

```
/*** print all elements in a tree by traversing inorder ***/
int print(BINARY_TREE_TYPE tree) {
   printf("Contents of tree by inorder traversal: \n");
   inorder(tree,0);
   printf("--- \n");
   return(0);
}
```

```
/*** error handler:
    print message passed as argument and take appropriate action ***/
int error(char *s) {
    printf("Error: %s\n",s);
    exit(0);
}
```

```
/*** assign values to an element ***/
int assign_element_values(ELEMENT_TYPE *e, int number, char s[]) {
    e->string = (char *) malloc(sizeof(char) * (strlen(s)+1));
    strcpy(e->string, s);
    e->number = number;
    return(0);
}
```

```
int main() {
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   BINARY TREE TYPE tree;
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   assign_element_values(&e, 3, "...");
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   assign_element_values(&e, 3, "...");
   delete element(e, &tree);
   print(tree);
```

