

04-630

# Data Structures and Algorithms for Engineers

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# Lecture 11

Queue ADT

Implementation using List ADT (array and linked-list)

Comparison of order of complexity

Dedicated ADT

Circular queues

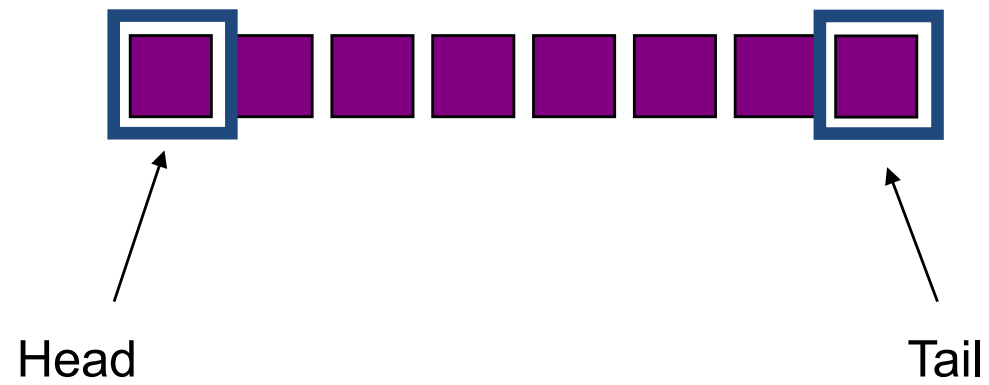
Queue applications

# Queues

A queue is another special type of list

- insertions are made at one end, called the tail of the queue
- deletions take place at the other end, called the head
- thus, the last one added is always the last one available for deletion
- also referred to as
  - FIFO list (First In First Out)

# Queues



# Queue Operations

*Declare*:  $\rightarrow Q$  :

The function value of *Declare*( $Q$ ) is an empty queue

# Queue Operations

*Empty.*  $\rightarrow Q$  :

The function *Empty* causes the queue to be emptied and it returns position *End(Q)*



# Queue Operations

*IsEmpty*.  $\mathbf{Q} \rightarrow \mathbf{B}$  :

The function value *IsEmpty*( $Q$ ) is *true* if  $Q$  is empty; otherwise it is *false*

# Queue Operations

*Head*:  $\mathbf{Q} \rightarrow \mathbf{E}$  :

The function value  $Head(Q)$  is the first element in the list;

if the queue is empty, the value is undefined

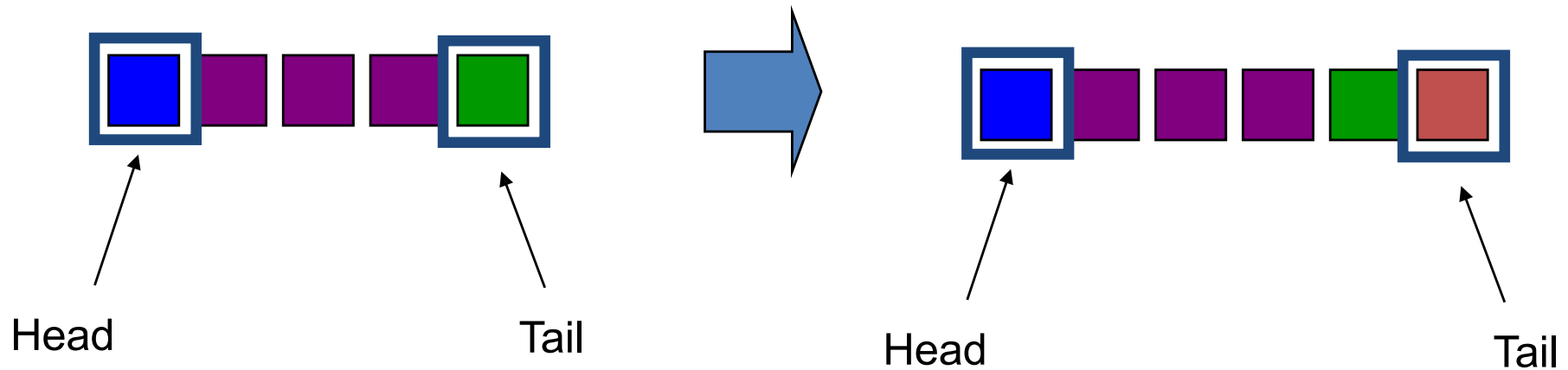


# Queue Operations

*Enqueue:*  $E \cup Q \rightarrow Q$  :

*Enqueue*( $e, Q$ )

Add an element  $e$  to the tail of the queue

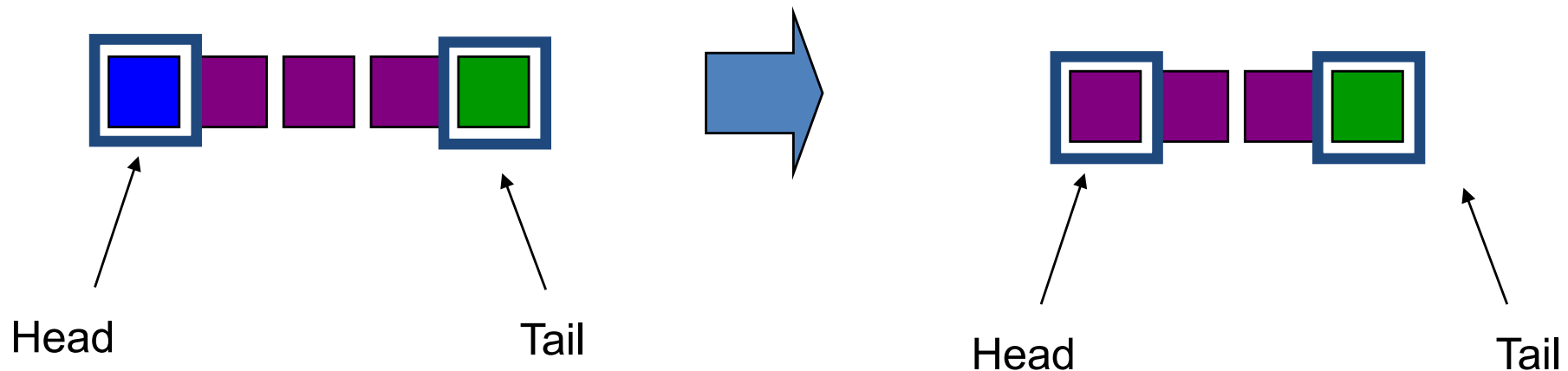


# Queue Operations

*Dequeue:* **Q** → **E** :

*Dequeue(Q)*

Remove the element from the head of the queue: i.e. return the first element and delete it from the queue



# Queue Operations

- All these operations can be directly implemented using the LIST ADT operations on a queue  $Q$
- Again, it may be more efficient to use a dedicated implementation
- And, again, it depends what you want: code efficiency or software re-use (i.e. utilization efficiency)

# Queue Operations

Declare(Q)

Empty(Q)

Head(Q)

Retrieve(First(Q), Q)

Enqueue(e, Q)

Insert(e, End(Q), Q)

Deque(Q)

Retrieve(First(Q), Q)

Delete(First(Q), Q)

# Queue Errors

- Queue **overflow** errors occur when you attempt to **enqueue()** an element in a queue that is **full**
- Queue **underflow** errors occur when you attempt to **dequeue()** an element from an empty queue
- Your ADT implementation should provide guards that catch these errors

# Queue Implementation

- The List ADT can be implemented
  - As an array
  - As a linked-list
- So, therefore, so can the Queue ADT
- What are the advantages and disadvantages of the these two options?
- When would you pick one implementation over the other?

# Queue Operations

	Array	Linked-List
Declare(Q)	$O(1)$	$O(1)$
Empty(Q)	$O(1)$	$O(n)$
Head(Q) Retrieve(First(Q), Q)	$O(1)$	$O(1)$
Enqueue(e, Q) Insert(e, End(Q), Q)	$O(1)$	$O(n)$ ... why?
Dequeue(Q) Retrieve(First(Q), Q) Delete(First(Q), Q)	$O(n)$ ... why?	$O(1)$

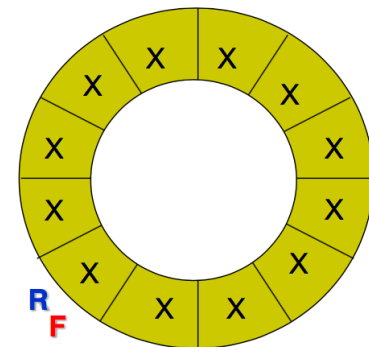
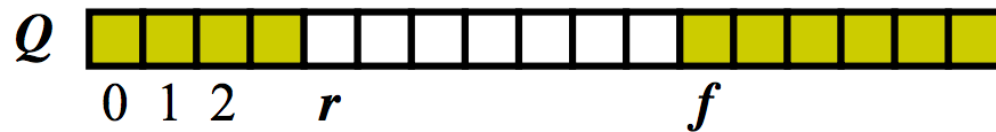
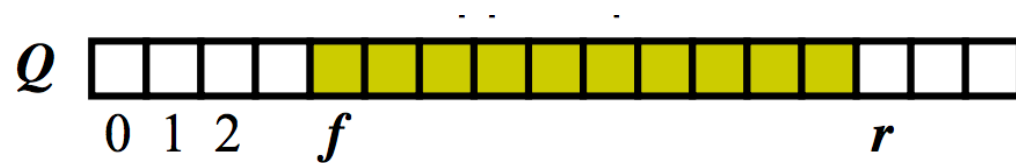
# Queue Implementation

- Reusing the List ADT involves some compromises
- Alternative is to create a new Queue ADT
  - With an implementation that avoids these compromises



# Queue Implementation

```
typedef struct {  
    int front;  
    int rear;  
    ITEM_TYPE items[MAX_ITEMS];  
} QUEUE_TYPE;
```



# Queue Implementation

```
void empty(Queue_TYPE *queue) {  
    queue->front = 0;  
    queue->rear = 0;  
    return(end(queue));  
}
```

# Queue Implementation

```
bool is_empty(Queue_TYPE *queue) {  
    if (queue->front == queue->rear)  
        return(true);  
    else  
        return(false)  
}
```

# Queue Implementation

```
void enqueue(ELEMENT_TYPE e, QUEUE_TYPE *queue) {  
    items[rear] = e;  
    rear = (rear + 1) % MAX_ITEMS;  
}
```

# Queue Implementation

```
int is_full(Queue_TYPE *queue) {  
    if (( rear + 1 ) % MAX_ITEMS == front )  
        return(TRUE);  
    else  
        return(FALSE);  
}
```

# Queue Implementation

```
void enqueue(ITEM_TYPE e, QUEUE_TYPE *queue) {
    if (!is_full(queue)) {
        items[rear] = e;
        rear = (rear + 1) % MAX_ITEMS;
    }
    else {
        error("Queue overflow: queue is already full");
    }
}
```

# Queue Implementation

```
void dequeue(ITEM_TYPE *e, QUEUE_TYPE *queue) {
    if (!is_empty(queue)) {
        *e = items[front];
        front = (front+1) % MAX_ITEMS;
    }
    else {
        error("Queue underflow: queue is empty");
    }
}
```

# Queue Implementation

- Can you see a particular problem with the linked-list implementation?
- How would you fix it?



# Queue Operations

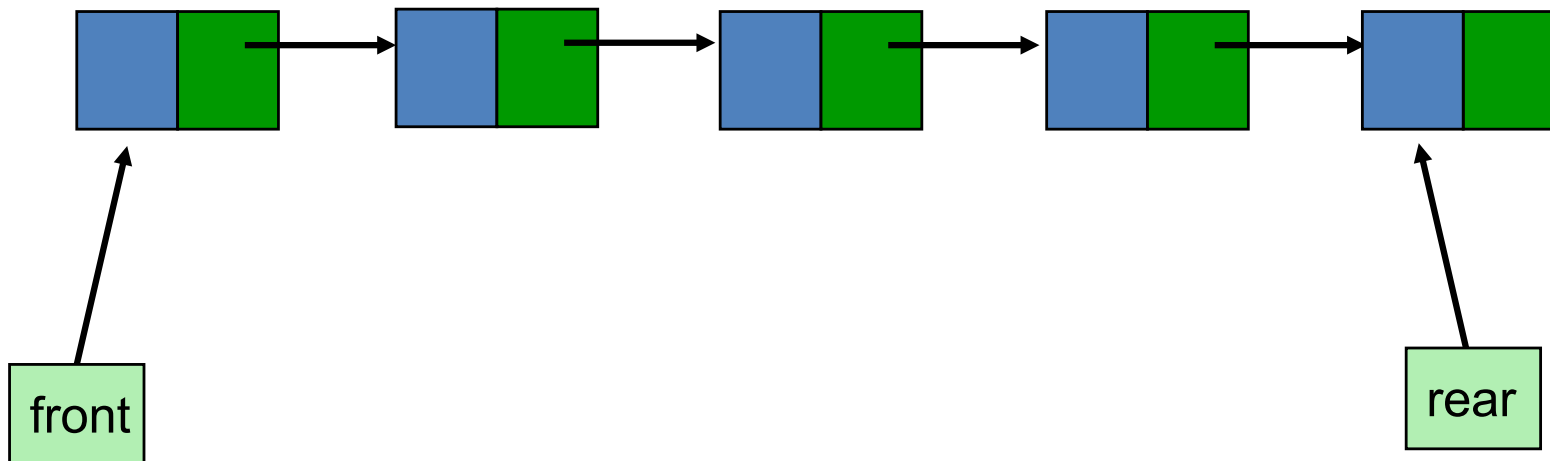
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Dequeue(Q) Retrieve(First(Q), Q) Delete(First(Q), Q)	$O(n)$ ... why?	$O(1)$

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# Queue Applications

- Scheduling/ waiting for system service queues
- Resource queues – provide coordinated access to shared resources
- Message queues (Buffer)
- Multi-queues and priority queues