04-630 Data Structures and Algorithms for Engineers

David Vernon
Carnegie Mellon University Africa

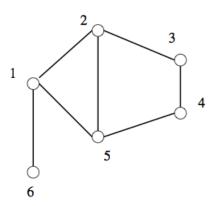
vernon@cmu.edu www.vernon.eu

Lecture 19

Graphs

- Types of graph
- Adjacency matrix representation
- Adjacency list representation
- Breadth-First Search (BFS) traversal
- Application of BFS
- Depth-First Search (DFS) traversal
- Topological Sorting
- Minimum Spanning Tree
 - Prim's Algorithm
 - Kruskal's algorithm
- Shortest Path Algorithms
 - Dijkstra's algorithm
 - Floyd's algorithm

- Visit every vertex and edge in a systematic way
- Key idea: mark each vertex when we first visit it & keep track of what we have not yet completely explored



- Each vertex will exist in one of three states
 - Undiscovered the vertex is in its initial untouched state
 - 2. Discovered the vertex has been found, but we have not yet processed all its edges
 - 3. Processed the vertex after we have visited all its edges

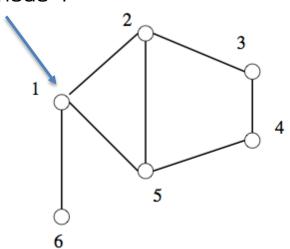
- Keep a record of all the vertices discovered but not yet completely processed
- Begin with a starting vertex
- Explore each vertex
 - Evaluate each edge leaving it
 - If the edge goes to an undiscovered vertex
 - Mark it discovered
 - Add it to the list of work to do
 - If the edge goes to a processed vertex, ignore it
 - If the edge goes to a discovered unprocessed vertex, ignore it

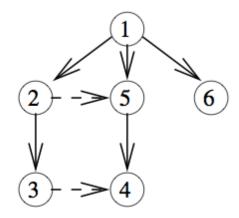
- There are two primary graph traversal algorithms
 - Breadth-first search (BFS)
 - Depth-first search (DFS)
- The difference is the order in which they explore vertices

- The order depends completely on the container data structure used to store the discovered but not processed vertices
 - BFS uses a queue
 - By storing the vertices in a FIFO queue, we explore the oldest unexplored vertices first
 - Thus explorations radiate out slowly from the starting vertex
 - DFS uses a stack
 - By storing the vertices in a LIFO stack, we explore the vertices by diving down a path, visiting a new neighbour if one is available, and backing up only when we are surrounded by (i.e. connected by edges to) previously discovered vertices
 - Thus explorations quickly wander away from out starting vertex

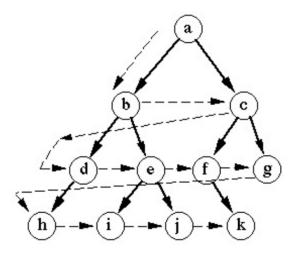
Breadth-first search (BFS)

Start at node 1

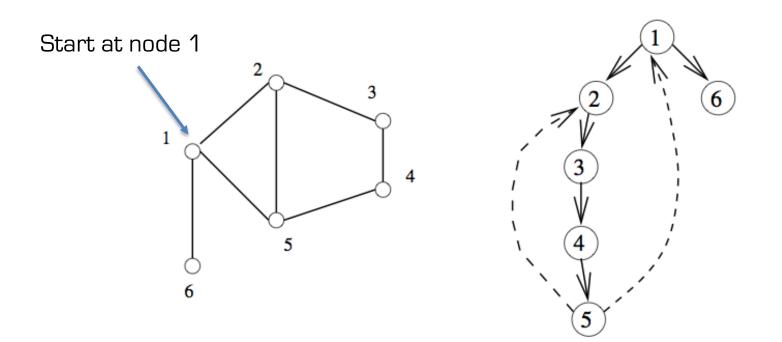




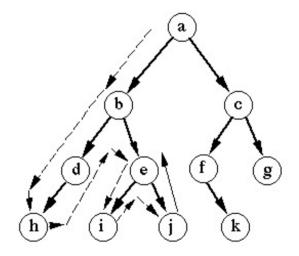
Breadth-first search (BFS)



Depth-first search (DFS)



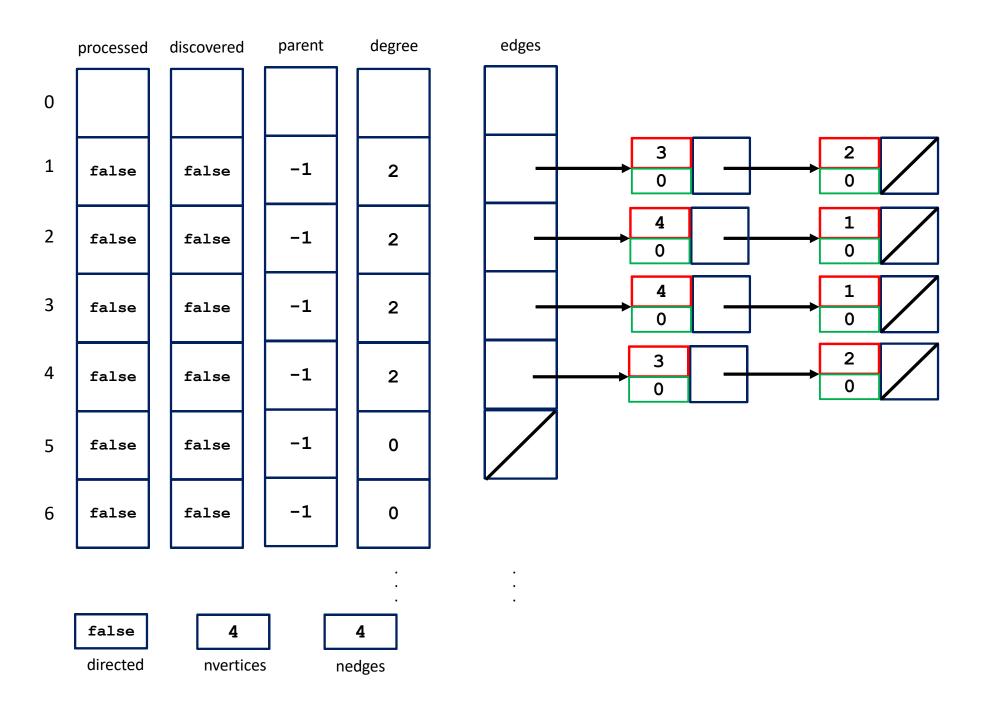
Depth-first search (DFS)



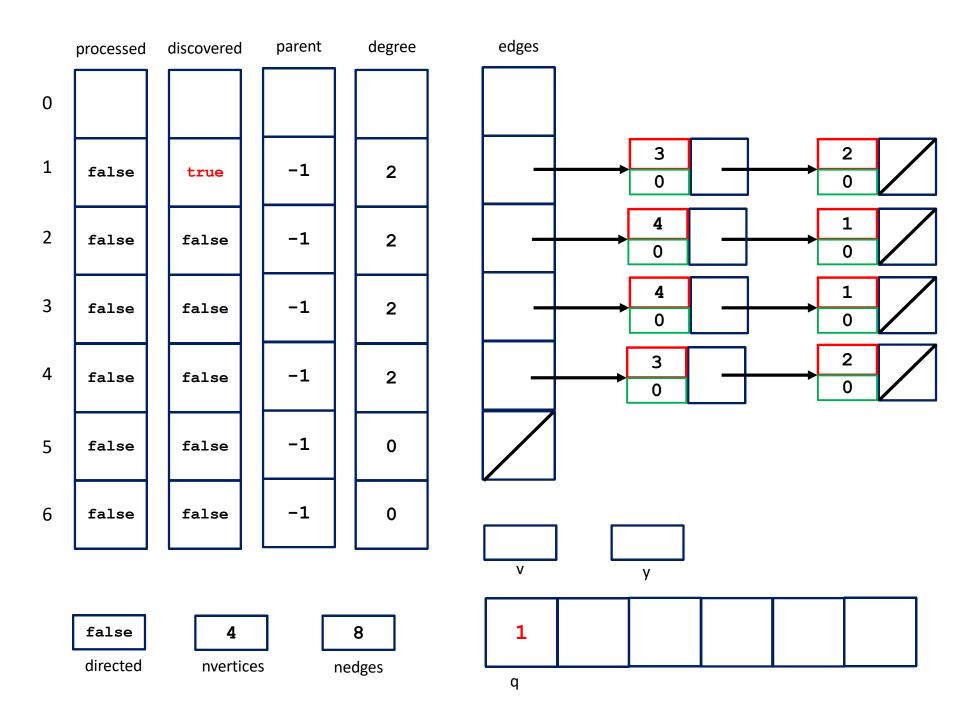
- Assign a direction to each edge, from discoverer vertex u to discovered vertex v
- Since each node has exactly one parent, except for the root (i.e. start vertex), this defines a tree on the vertices of the graph
- This tree defines the shortest path from the root to every other node in the tree
- This makes the BFS very useful for in shortest path problems (in unweighted graphs)

```
\mathrm{BFS}(G,s)
      for each vertex u \in V[G] - \{s\} do
             state[u] = "undiscovered"
            p[u] = nil, i.e. no parent is in the BFS tree
      state[s] = "discovered"
      p[s] = nil
      Q = \{s\}
      while Q \neq \emptyset do
             u = \text{dequeue}[Q]
             process vertex u as desired
             for each v \in Adj[u] do
                   process edge (u, v) as desired
                   if state[v] = "undiscovered" then
                          state[v] =  "discovered"
                          p[v] = u
                          enqueue[Q, v]
             state[u] = "processed"
```

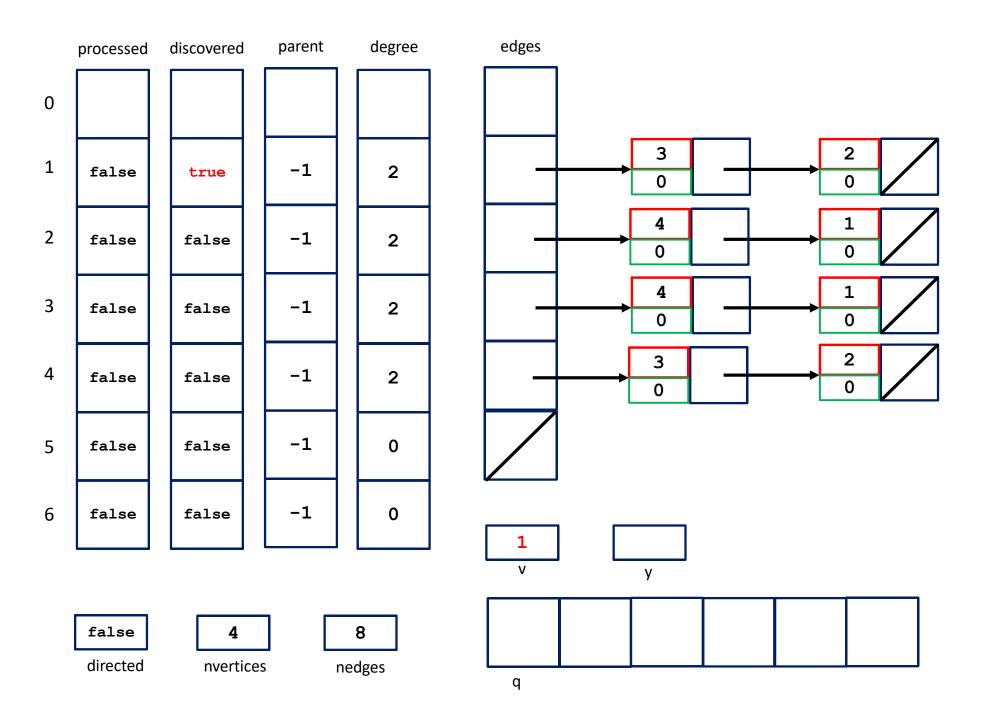
```
/* Breadth-First Search
                                                                    */
bool processed [MAXV+1]; /* which vertices have been processed */
bool discovered[MAXV+1]; /* which vertices have been found */
int parent[MAXV+1]; /* discovery relation */
/* Each vertex is initialized as undiscovered:
                                                                    */
initialize search(graph *g) {
                                   /* counter */
   int i;
   for (i=1; i<=q->nvertices; i++) {
      processed[i] = discovered[i] = false;
     parent[i] = -1;
```



```
/* Once a vertex is discovered, it is placed on a queue.
                                                                    */
/* Since we process these vertices in first-in, first-out order,
                                                                    */
/* the oldest vertices are expanded first, which are exactly those */
/* closest to the root
                                                                    */
bfs(graph *g, int start)
                             /* queue of vertices to visit */
   queue q;
   int v;
                             /* current vertex
                                                            */
                            /* successor vertex
                                                            */
   int y;
                                                            */
                            /* temporary pointer
   edgenode *p;
   init queue(&q);
   enqueue(&q,start);
   discovered[start] = true;
```

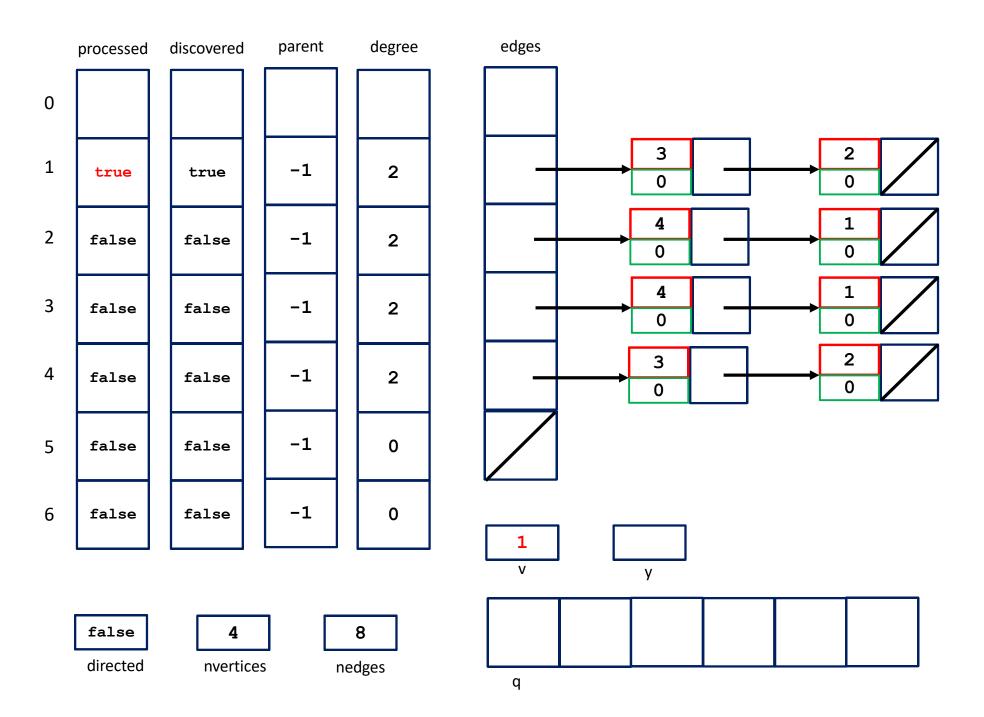


```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = TRUE;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == FALSE) || g->directed)
         process_edge(v,y);
      if (discovered[y] == FALSE) {
         enqueue(&q,y);
         discovered[y] = TRUE;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```

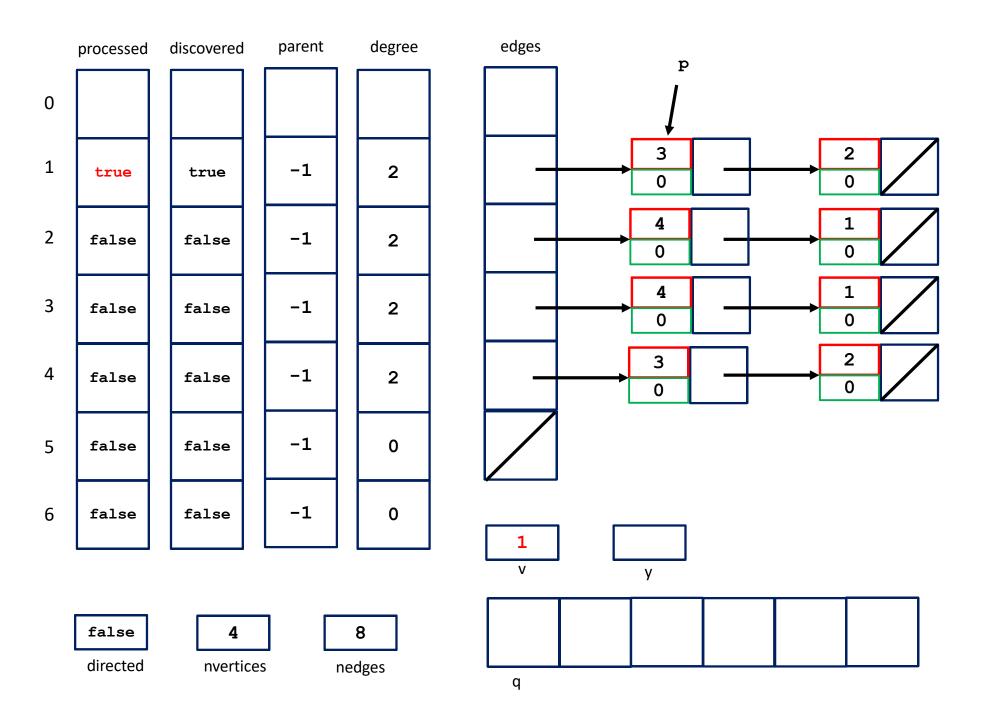


```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = TRUE;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == FALSE) || g->directed)
         process_edge(v,y);
      if (discovered[y] == FALSE) {
         enqueue(&q,y);
         discovered[y] = TRUE;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```

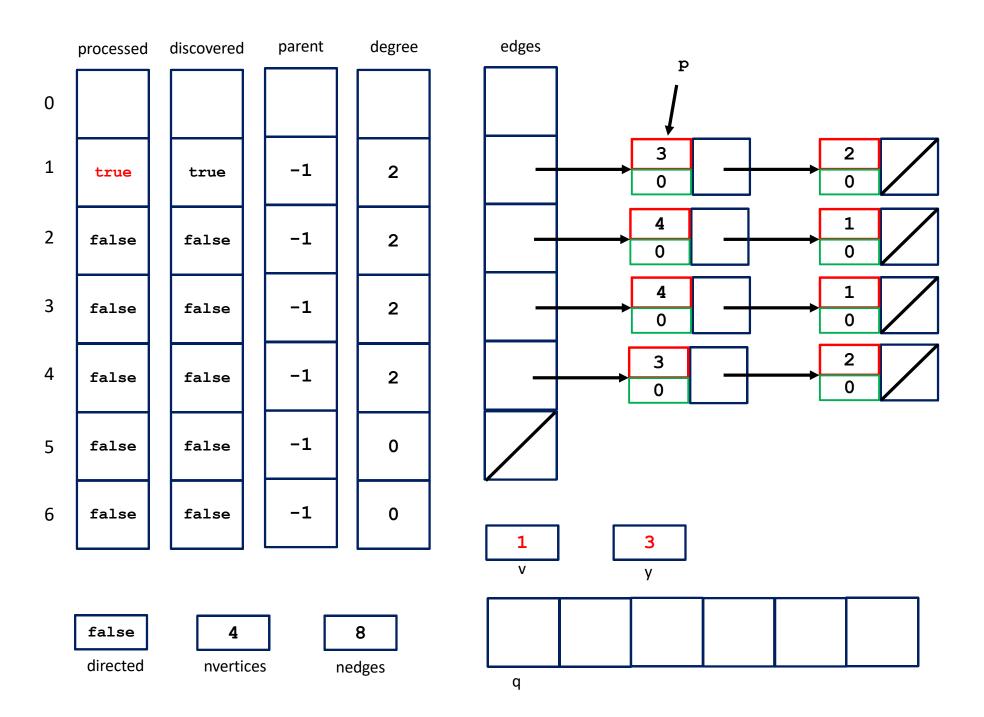
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = TRUE;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == FALSE) || g->directed)
         process_edge(v,y);
      if (discovered[y] == FALSE) {
         enqueue(&q,y);
         discovered[y] = TRUE;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```



```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = TRUE;
   p = g->edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == FALSE) || g->directed)
         process_edge(v,y);
      if (discovered[y] == FALSE) {
         enqueue(&q,y);
         discovered[y] = TRUE;
         parent[y] = v;
      p = p->next;
  process vertex late(v);
```

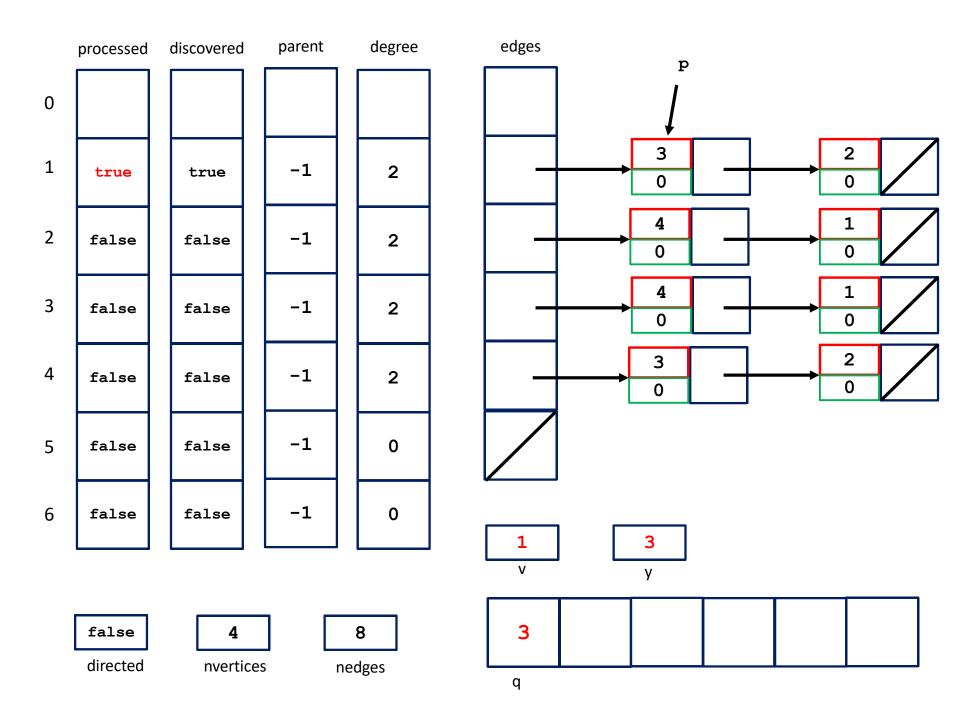


```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = TRUE;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == FALSE) || g->directed)
         process_edge(v,y);
      if (discovered[y] == FALSE) {
         enqueue(&q,y);
         discovered[y] = TRUE;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```

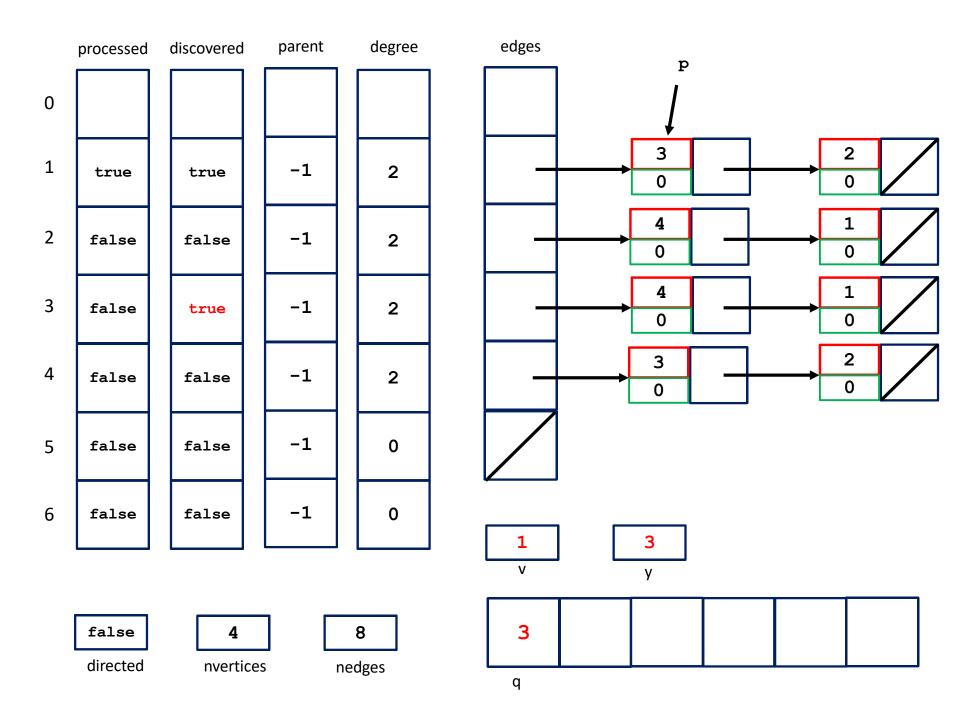


```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```

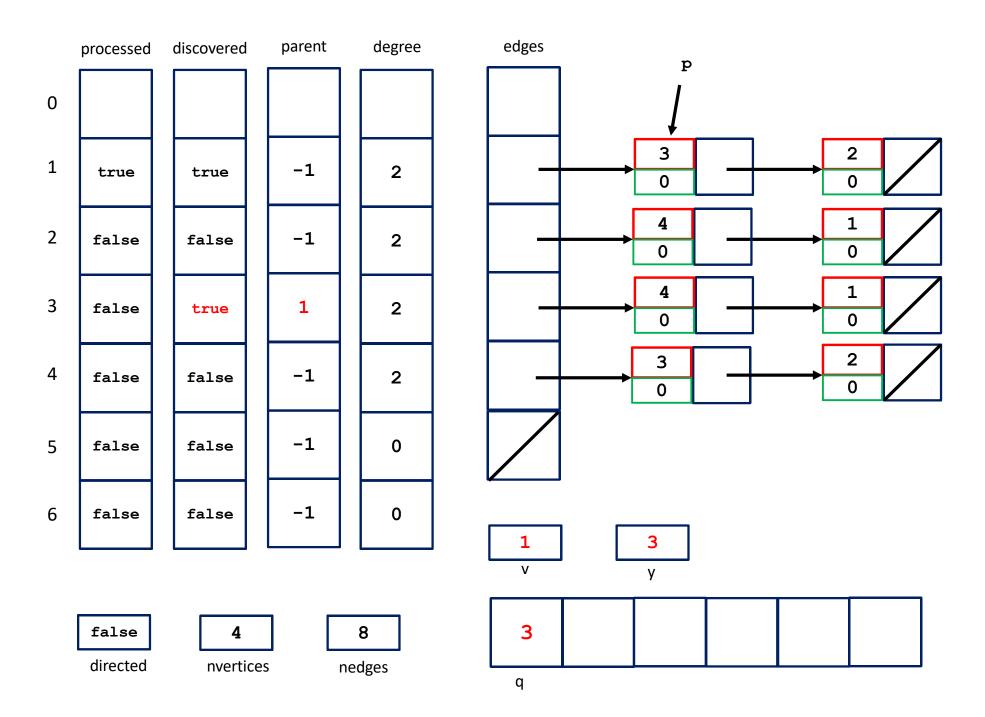
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```



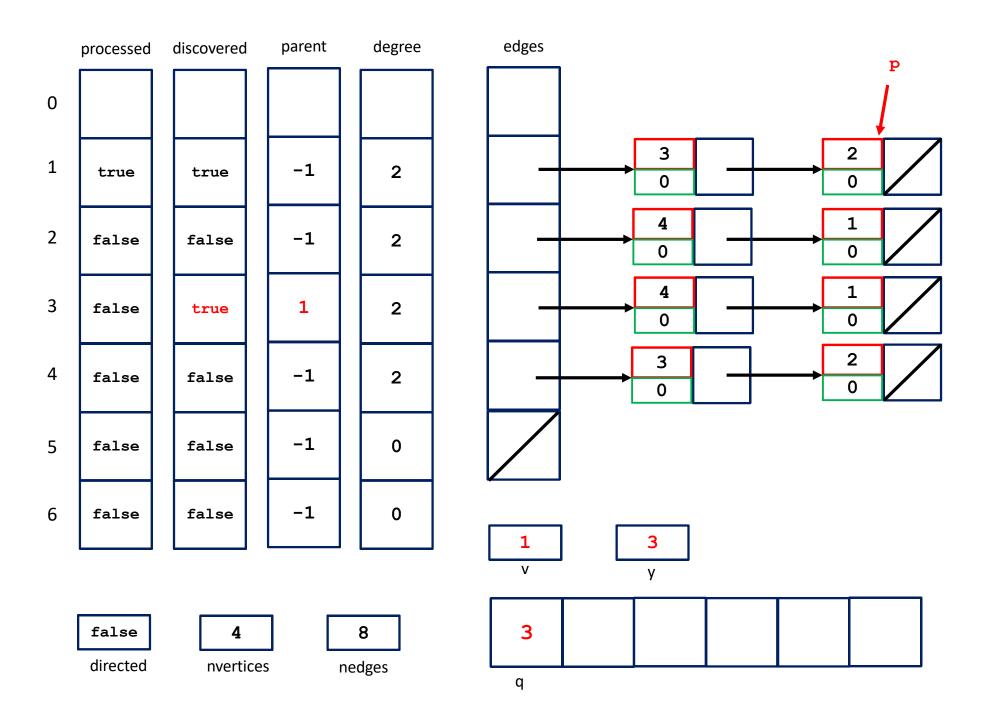
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```



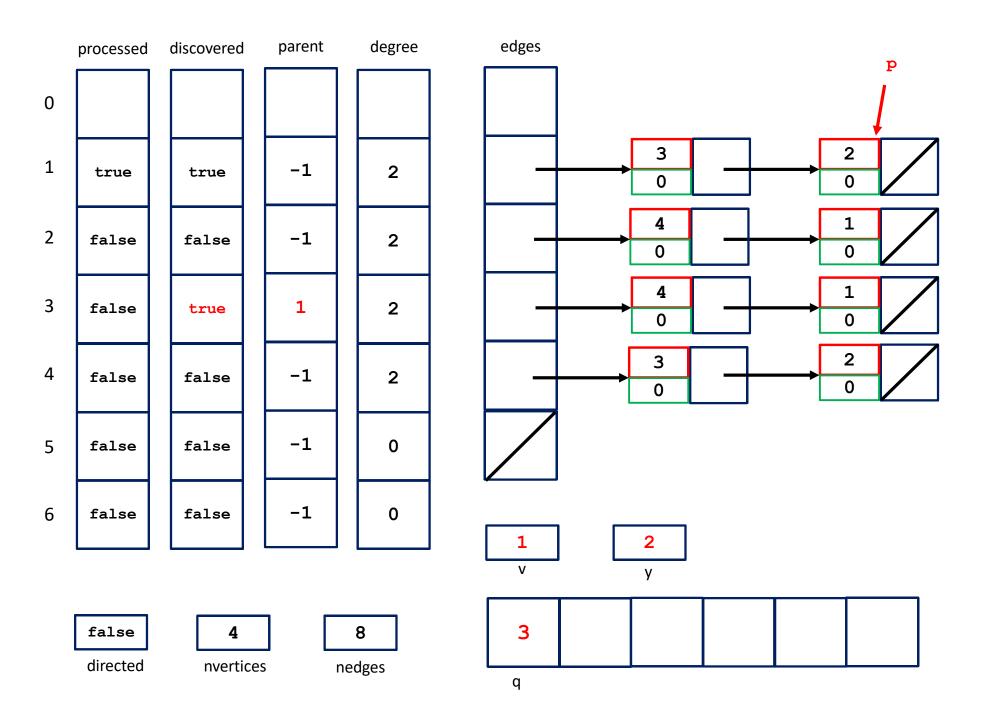
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```



```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```

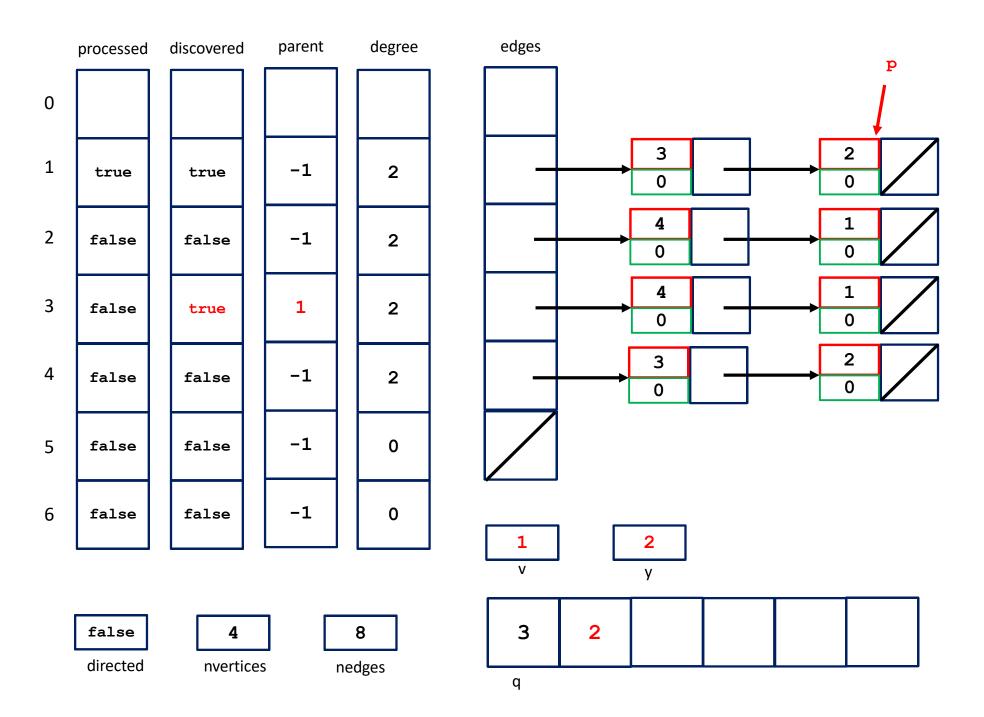


```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = TRUE;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == FALSE) || g->directed)
         process_edge(v,y);
      if (discovered[y] == FALSE) {
         enqueue(&q,y);
         discovered[y] = TRUE;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```

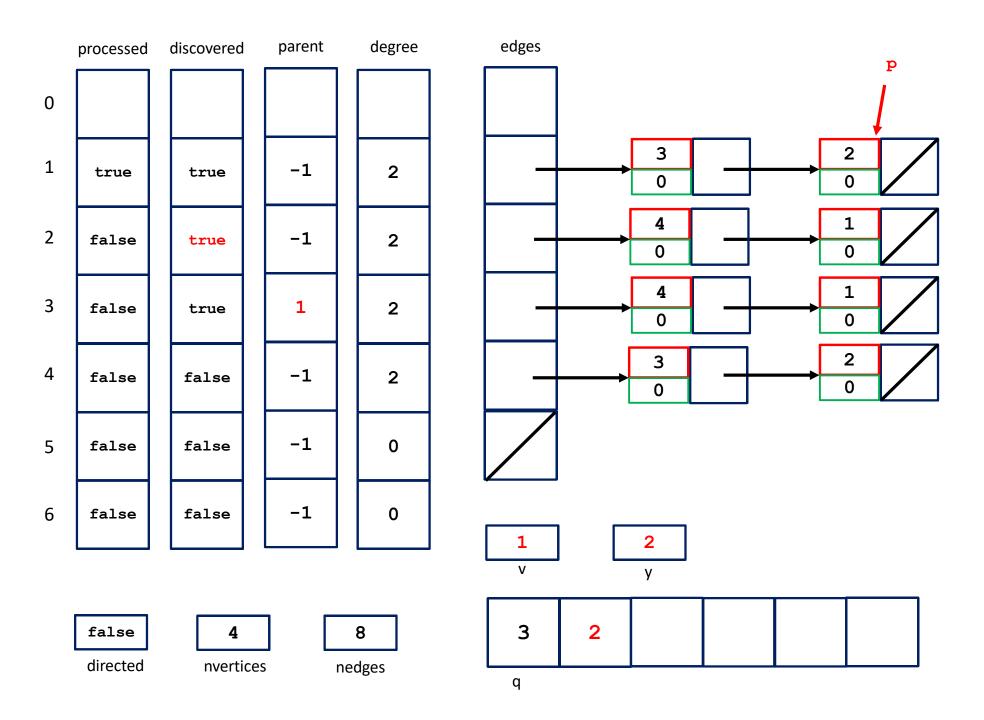


```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```

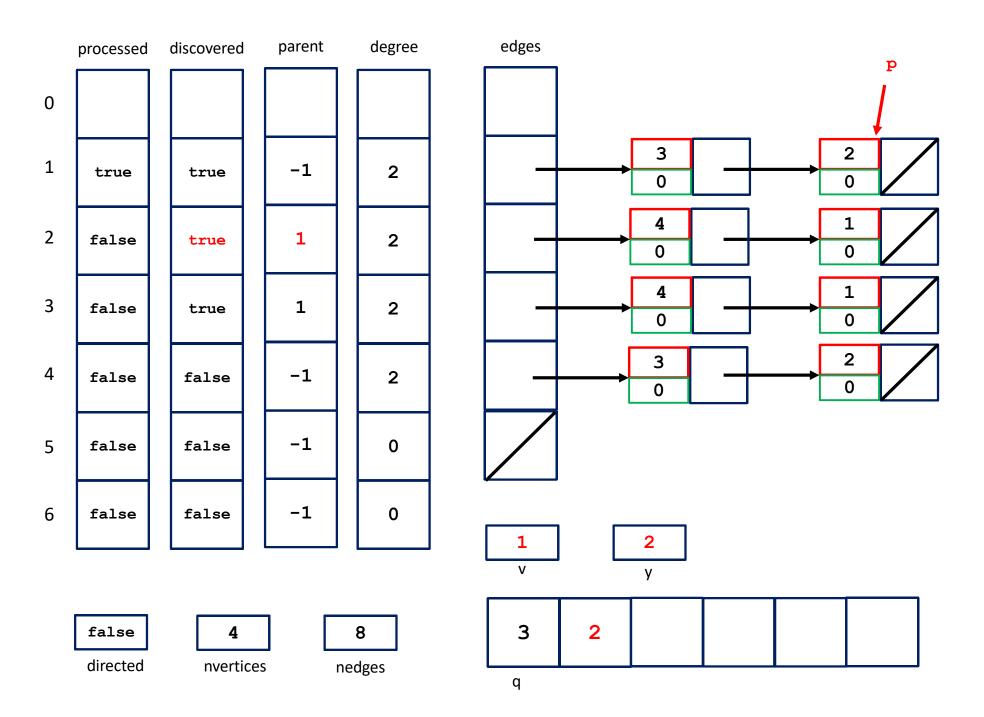
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```



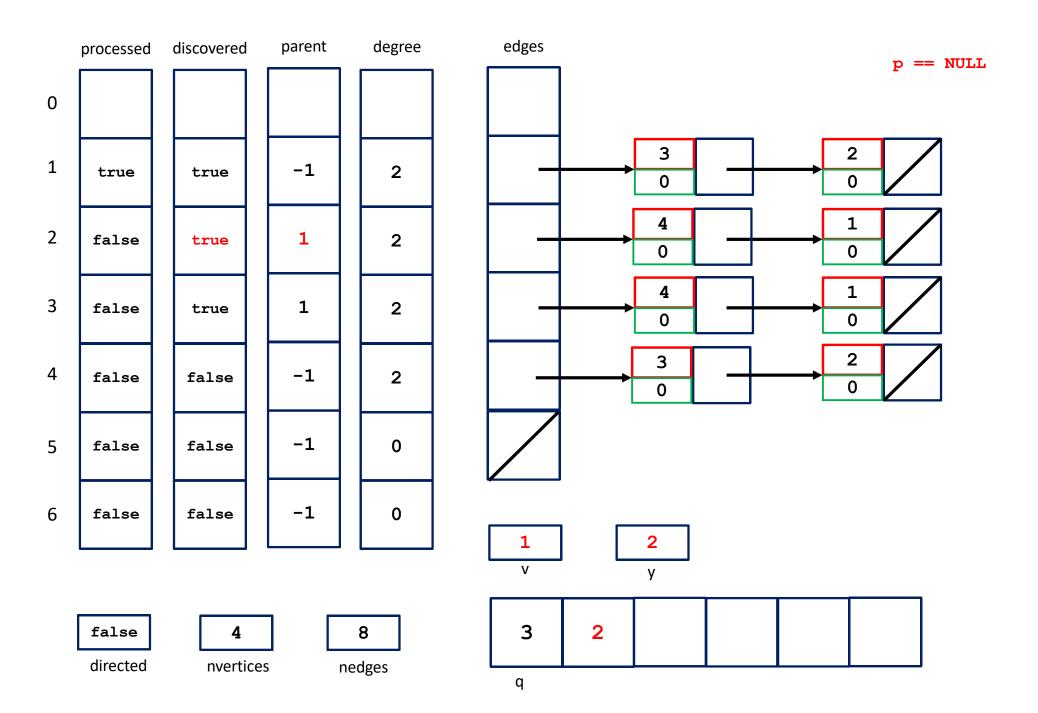
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```



```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```

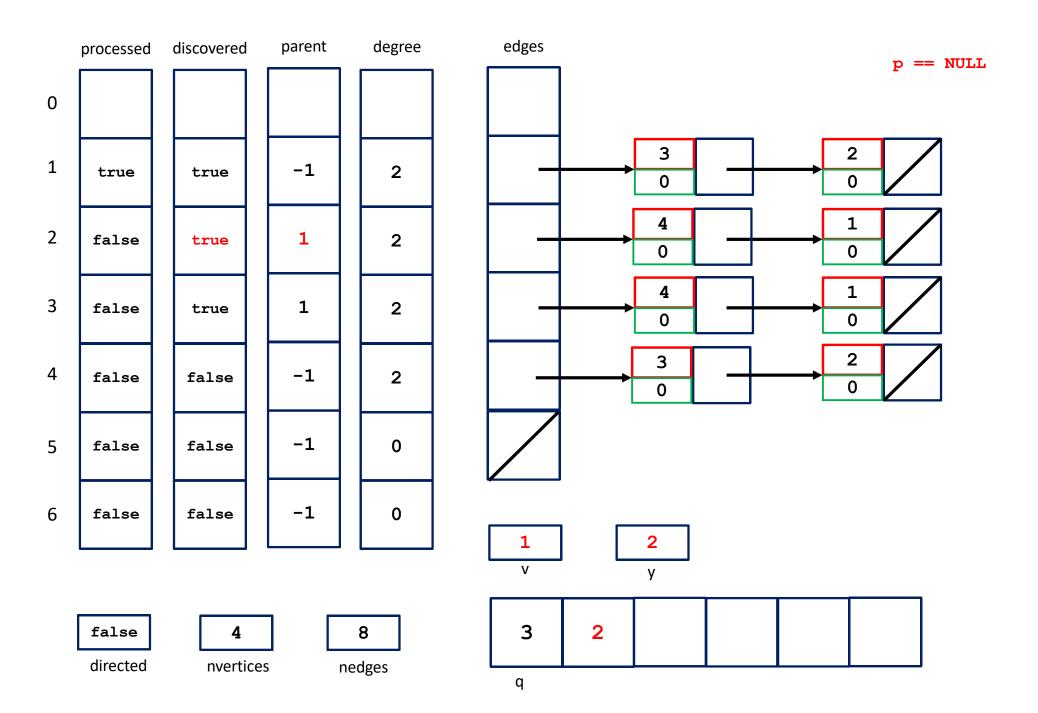


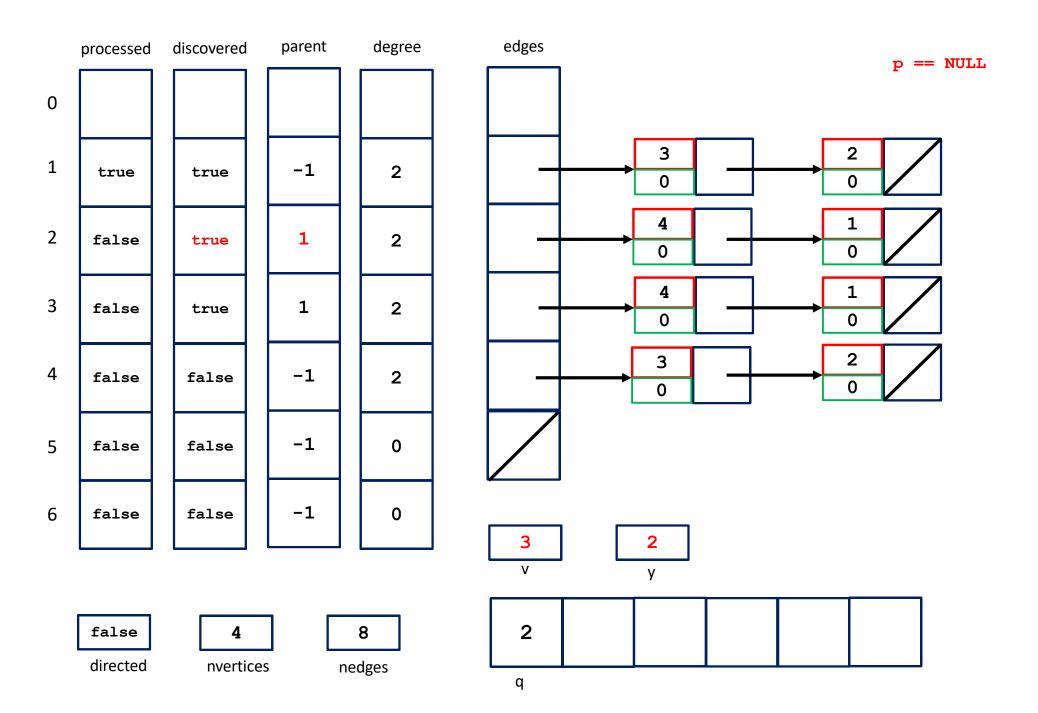
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process_edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```



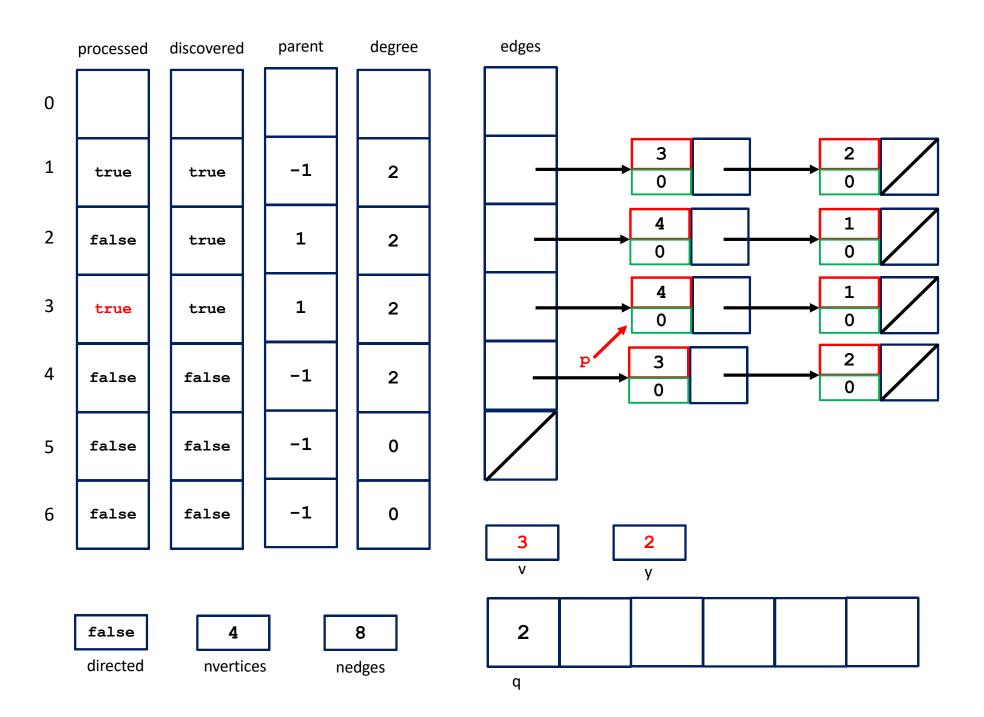
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex_late(v);
```

```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```





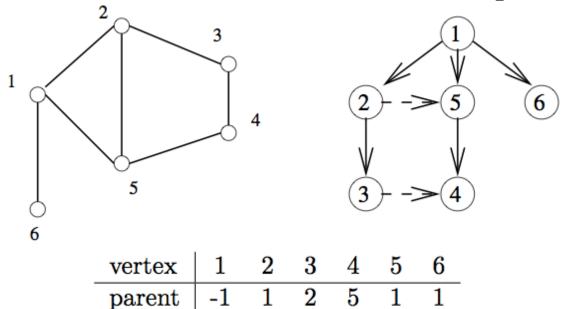
```
while (empty queue(&q) == FALSE) {
   v = dequeue(&q);
   process vertex early(v);
   processed[v] = true;
   p = g \rightarrow edges[v];
   while (p != NULL) {
      y = p->y;
      if ((processed[y] == false) || g->directed)
         process edge(v,y);
      if (discovered[y] == false) {
         enqueue(&q,y);
         discovered[y] = true;
         parent[y] = v;
      p = p->next;
   process vertex late(v);
```



```
/* The exact behaviour of bfs depends on the functions
                                                                     */
                                                                     */
/*
      process vertex early()
                                                                     */
/*
     process vertex late()
                                                                     */
      process edge()
/*
/* These functions allow us to customize what the traversal does
                                                                     */
                                                                     */
/* as it makes its official visit to each edge and each vertex.
/* Here, e.g., we will do all of vertex processing on entry
                                                                     */
                                                                     */
/* (to print each vertex and edge exactly once)
                                                                     */
/* so process vertex late() returns without action
process vertex late(int v) {
process vertex early(int v) {
   printf("processed vertex %d\n",v);
process edge(int x, int y) {
   printf("processed edge (%d,%d)\n",x,y);
}
```

Finding Paths

- The parent array in bfs() is very useful for finding interesting paths through a graph
- The vertex that discovered vertex i is defined as parent[i]



Finding Paths

- Every vertex is discovered during the course of a traversal so every node has a parent (except the root)
- The parent relation defines a tree of discovery with the initial search node as the root of the tree
- Because vertices are discovered in order of increasing distance from the root, this tree has a very important property
 - The unique tree path from the root to each node uses the smallest number of edges (and intermediate nodes) possible on any path from the root to that vertex
 - Thus BFS can be used to find shortest paths in an unweighted graph

Finding Paths

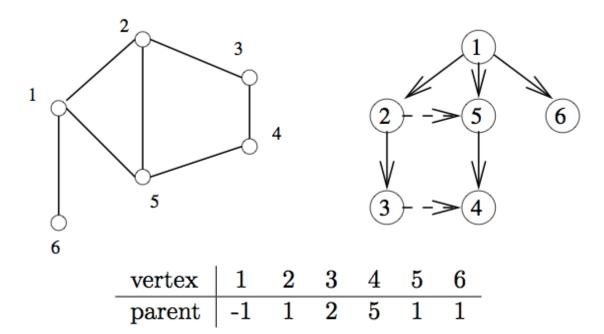
- To reconstruct a path we follow the chain of ancestors from the destination node x to the root
- Note we have to work backwards (we only know the parents)
- We find the path from to the root and
 - Either store it and explicitly reverse it using a stack
 - Or construct the path recursively (in which case the stack is implicit)

```
bool find path(int start, int end, int parents[]) {
    bool is path;
    if (end == -1) {
        is path = false; // some vertex on the path back from the end
                                // has no parent (not counting start)
    else if ((start == end)) {
         printf("\n%d", start);
         is path = true; // we have reached the start vertex
    else {
        is path = find path(start,parents[end],parents);
        printf(" %d",end);
    return(is path);

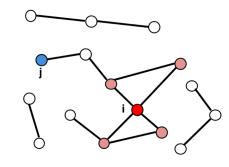
        vertex
        1
        2
        3
        4
        5
        6

        parent
        -1
        1
        2
        5
        1
        1
```

find_path(1,4,parent)

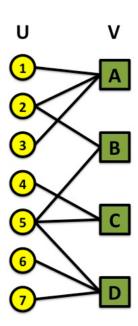


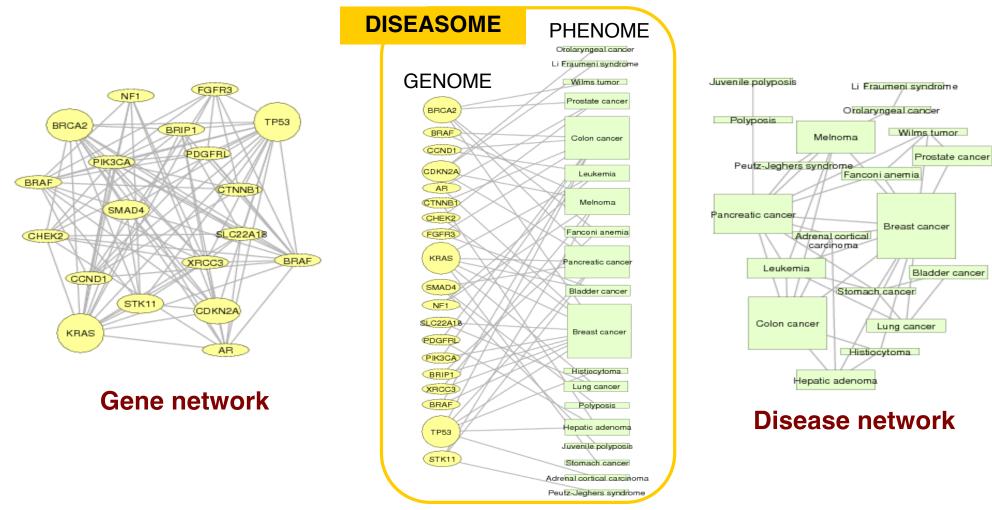
Applications of Breadth-First Search



- Identifying connected components
 - A graph is connected if there is a path between any two vertices
 - A connected component of an undirected graph is a maximal set of vertices such that there is a path between every pair of vertices
 - The components are separate "pieces" of the graph such that there is no connection between the pieces
 - Many complicated problems reduce to finding or counting connected components
 - How would you find and label all the components in a graph?

- Applications of Breadth-First Search
 - Two-Colouring Graphs
 - The vertex-colouring problem seeks to assign a label (or colour) to each vertex of a graph such that no edge links any two vertices of the same colour
 - The goal is use as few colours as possible
 - A graph is bipartite if it can be coloured without conflicts using only two colours



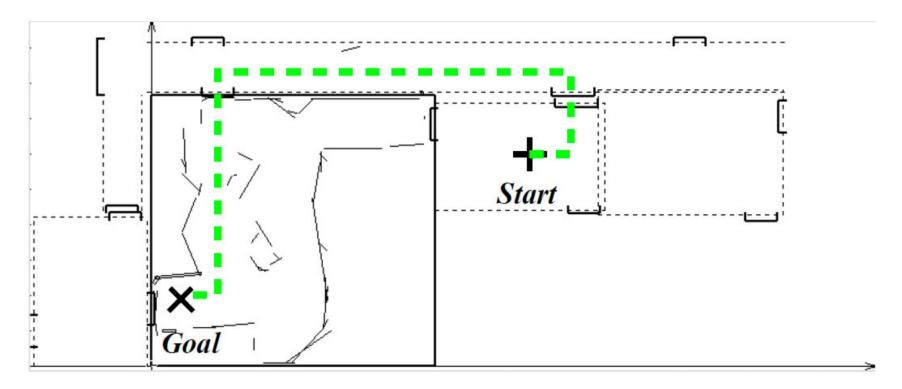


Goh, Cusick, Valle, Childs, Vidal & Barabási, PNAS (2007)

- Applications of Breadth-First Search
 - Robot path-planning

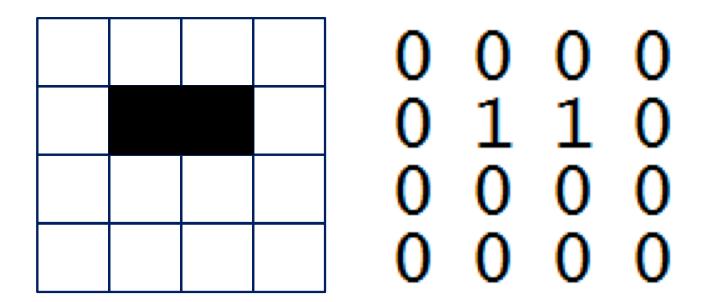


- Applications of Breadth-First Search
 - Robot path-planning



- Applications of Breadth-First Search
 - Robot path-planning

Represent the map of the environment as an occupancy grid



- Applications of Breadth-First Search
 - Robot path-planning

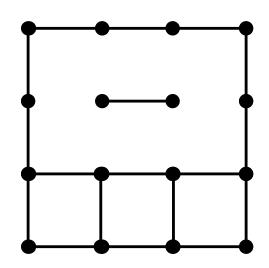
Represent the map of the environment as an occupancy grid

0	0	0	0
0	1	1	0
0	0	0	0
0	0	0	0

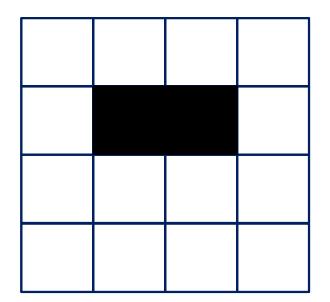
0	0	0	0
0	1	1	0
0	0	0	0
0	0	0	0

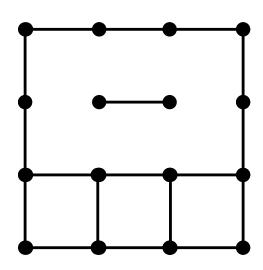
- Applications of Breadth-First Search
 - Robot path-planning
 Convert this to a graph

0	0	0	0
0	1	1	0
0	0	0	0
0	0	0	0



- Applications of Breadth-First Search
 - Robot path-planning
 Convert this to a graph

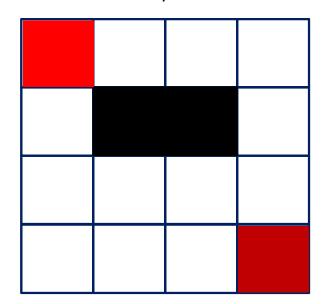


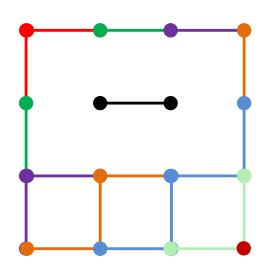


- Applications of Breadth-First Search
 - Robot path-planning

Do a BFS from the robot start position ...

To find the shortest path to all other vertices

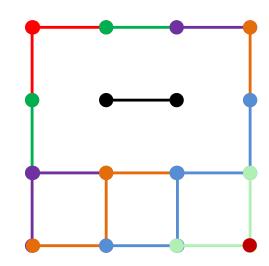




- Applications of Breadth-First Search
 - Robot path-planning

Mark the path from the robot start position to the goal position on the occupancy grid

2	0	0	0
2	1	1	0
2	0	0	0
2	2	2	2



Applications of Breadth-First Search

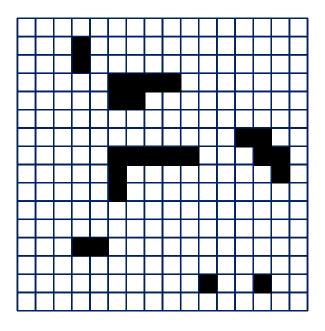
Robot path-planning

Mark the path from the robot start position to the goal position on the occupancy grid

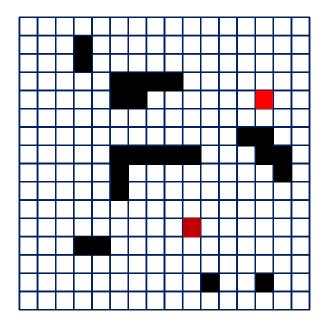
2	0	0	0
2	1	1	0
2	0	0	0
2	2	2	2

2	0	0	0
2	1	1	0
2	0	0	0
2	2	2	2

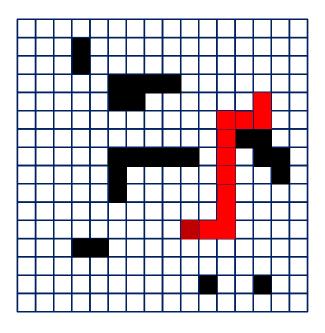
- Applications of Breadth-First Search
 - Robot path-planning



- Applications of Breadth-First Search
 - Robot path-planning



- Applications of Breadth-First Search
 - Robot path-planning



- Applications of Breadth-First Search
 - Robot path-planning

